## T-79.5501 Cryptology <br> Spring 2009

Homework 8

Tutor: Joo Y. Cho<br>joo.cho@tkk.fi

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Q1. Test the primality of 2009 using

1. the Solovay-Strassen test with $a=442$
2. the Miller-Rabin test with $a=442$.

A1-a). (The Solovay-Strassen test)
We proceed as written in the lecture slides. We have $n=2009$, $a=442=2 \cdot 13 \cdot 17$.

$$
\begin{aligned}
& \left(\frac{442}{2009}\right)=\left(\frac{2}{2009}\right)\left(\frac{13}{2009}\right)\left(\frac{17}{2009}\right)=\left(\frac{13}{2009}\right)\left(\frac{17}{2009}\right) \\
& =\left(\frac{7}{13}\right)\left(\frac{3}{17}\right)=\left(\frac{6}{7}\right)\left(\frac{2}{3}\right)=-\left(\frac{3}{7}\right)=1
\end{aligned}
$$

Also, we compute $442^{\frac{n-1}{2}}=442^{1004} \bmod 2009$. Since $442^{4} \equiv 1 \bmod 2009$, we get $442^{1004}=442^{4 \cdot 251} \equiv 1 \bmod 2009$.
Hence, $\left(\frac{a}{n}\right) \equiv a^{\frac{n-1}{2}}$ so $n$ is prime.

A1-b). (The Miller-Rabin test) We have $n=2009, a=442$, $n-1=2^{3} \cdot 251$ and $k=3$.

- $b \leftarrow 442^{251} \bmod 2009=50$ by square-and-multiply.
- $b \not \equiv 1 \bmod n$, continue.
- $i=0: b \not \equiv-1(\bmod n), b \leftarrow b^{2}=491$.
- $i=1: b \not \equiv-1(\bmod n), b \leftarrow b^{2}=1$.
- $i=2: b \not \equiv-1(\bmod n), b \leftarrow b^{2}=1$.
- Answer: " $n$ is composite".

Q2.

1. Find all square roots of 1 modulo $4453=61 \cdot 73$.
2. 2777 is a square root of 3586 modulo 4453 . Find all square roots of 3586 modulo 4453.

A2-a). The task is to find $x$ such that $x^{2} \equiv 1 \bmod 4453$.

- It is obvious that $x= \pm 1$ are square roots.
- Also, we have

$$
\begin{aligned}
x & \equiv-1(\bmod 61) \\
x & \equiv 1(\bmod 73)
\end{aligned}
$$

We get $61^{-1} \equiv 6 \bmod 73$ and $73^{-1} \equiv 56 \bmod 61$.

- Using CRT, we get $x=-1 \cdot 73 \cdot 56+1 \cdot 61 \cdot 6 \equiv 731 \bmod 4453$.
- In a similar way, we can get $x=-731$. Hence, $\pm 1$ and $\pm 731$ are four square roots of $1 \bmod 4453$.

A2-b). From the lecture slides,

- $\pm 1$ and $\pm 731$ are the square roots of $1 \bmod n$. Put $w=731$.
- Given a square root $b$ of $a$, the four square roots of $a \bmod n$ are $\pm b$ and $\pm b w$.
- So with $a=3586, b=2777, n=4453, w=731$, the four square roots of $3586 \bmod 4453$ are $\pm 2777$ and $\pm 2777 \cdot 731$, namely, $\{2777,1676,3872,581\}$.

Q3. (Stinson 5.24) Suppose that $i \geq 2$ and $b^{2} \equiv a\left(\bmod p^{i-1}\right)$. it was shown that there is a unique $x \in \mathbf{Z}_{p^{i}}$, such that $x^{2} \equiv a\left(\bmod p^{i}\right)$ and $x \equiv b\left(\bmod p^{i-1}\right)$ and

$$
\begin{align*}
b^{2} & =a+m p^{i-1} \bmod p^{i}  \tag{1}\\
x & =b+n p^{i-1} \bmod p^{i}  \tag{2}\\
n & =\frac{p-1}{2} b^{-1} m \bmod p \tag{3}
\end{align*}
$$

Starting with the congruence $6^{2} \equiv 17(\bmod 19)$, find square roots of 17 modulo $19^{2}$.

A3. We have $b^{2} \equiv a\left(\bmod p^{i-1}\right), x^{2} \equiv a\left(\bmod p^{i}\right)$ and $x \equiv b\left(\bmod p^{i-1}\right)$.

$$
\begin{aligned}
b^{2} & =a+m p^{i-1} \bmod p^{i} \\
x & =b+n p^{i-1} \bmod p^{i} \\
n & =\frac{p-1}{2} b^{-1} m \bmod p
\end{aligned}
$$

Using these equations, we find square roots of 17 modulo $19^{2}$ and modulo $19^{3}$.

1. $a=17, b=6, p=19$ and $i=2$.
2. $\mathrm{By}(1), b^{2}=36=17+1 \cdot 19$. We get $m=1$
3. $b^{-1} \bmod 19=16$. By (3), $n=9 \cdot 16 \cdot 1 \bmod 19=11$.
4. $x=6+11 \cdot 19 \bmod 19^{2}=215$ from (2).
5. Similarly, for $b=-6=13$, we get $x=146$.

A3. Find square roots of 17 modulo $19^{3}$.

$$
\begin{aligned}
b^{2} & =a+m p^{i-1} \bmod p^{i} \\
x & =b+n p^{i-1} \bmod p^{i} \\
n & =\frac{p-1}{2} b^{-1} m \bmod p
\end{aligned}
$$

Let now $i=3$.

1. $a=17, b=215, p=19$ and $i=3, p^{2}=361, p^{3}=6859$.
2. $\mathrm{By}(1), b^{2} \equiv 5071=17+14 \cdot 361$. We get $m=14$
3. $b^{-1} \bmod 19=16$. By (3), $n=9 \cdot 16 \cdot 14 \bmod 19=2$.
4. $x=215+2 \cdot 19^{2} \bmod 19^{3}=937$ from (2).

Similarly, for $b=-215 \bmod p^{2}=146$, we get $x=-937=5922$.

Q4. Compute

$$
2^{120}(\bmod 122183)
$$

Then using the $p-1$ method, attempt to factor 122183 .

A4.

- We calculate $2^{120} \bmod 122183$ by square-and-multiply as follows:

$$
2^{120}=2^{64} 2^{32} 2^{16} 2^{8} \equiv 15068 \bmod 122183
$$

- We also know that $120=5!=5 \cdot 4 \cdot 3 \cdot 2$.
- According to Pollard $p-1$, we set $a=2^{B!} \equiv 15068 \bmod 122183$ where $B=5$.
- Then, we calculate $d=\operatorname{gcd}(a-1, n)=\operatorname{gcd}(15067,122183)=61$. Since $1<d<n$, we conclude that 61 is a factor of 122183 . Indeed, we can see $122183=61 \cdot 2003$.
- Note this worked because all prime power divisors of $d-1=60=2^{2} \cdot 3 \cdot 5$ were less than or equal to $B=5$.

Q5. Let $n=p q$, where $p$ and $q$ are primes. We can assume that $p>q>2$ and we denote $d=\frac{p-q}{2}$ and $x=\frac{p+q}{2}$. Then $n=x^{2}-d^{2}$. Attempt to factor $n=400219845261001$ by searching for small non-negative integers $t$ such that $x^{2}-n=(\lceil\sqrt{n}\rceil+t)^{2}-n$ is a perfect square. (This is a simple form of the Quadratic Sieve method.)

A5. The task is to search for small non-negative $t$ such that $(\lceil\sqrt{n}\rceil+t)^{2}-n$ is a perfect square, and as a result we have an equation like $n=a^{2}-b^{2}=(a+b)(a-b)$ with $a, b$ known and we find the factors of $n$. We set $n=400219845261001$ and try for $t=1 \ldots$.

- $t \leftarrow 1,(20005496+1)^{2}-400219845261001=64956008$ is not a square.
- $t \leftarrow 2,(\lceil\sqrt{n}\rceil+2)^{2}-n=104967003$ is not a square.
- $t \leftarrow 3,(\lceil\sqrt{n}\rceil+3)^{2}-n=144978000$ is not a square.
- $t \leftarrow 4,(\lceil\sqrt{n}\rceil+4)^{2}-n=184988999$ is not a square.
- $t \leftarrow 5,(20005501)^{2}-n=225000000=15000^{2}$.

We now have the factors: $20005501 \pm 15000$ and $400219845261001=19990501 \times 20020501$. This worked because $p, q$ were too close to each other.

Q6.

1. Bob: $\left(n, b_{1}\right)$, Charlie : $\left(n, b_{2}\right)$, and $\operatorname{gcd}\left(b_{1}, b_{2}\right)=1$
2. Alice: $y_{1}=x^{b_{1}} \bmod n \Longrightarrow$ Bob and $y_{2}=x^{b_{2}} \bmod n \Longrightarrow$ Charlie
3. Oscar intercepts $y_{1}$ and $y_{2}$, and performs
i) Compute $c_{1}=b_{1}^{-1} \bmod b_{2}$
ii) Compute $c_{2}=\left(c_{1} b_{1}-1\right) / b_{2}$
iii) Compute $x_{1}=y_{1}^{c_{1}}\left(y_{2}^{c_{2}}\right)^{-1} \bmod n$
4. Prove that the value $x_{1}$ computed in step iii) is in fact Alice's plaintext, $x$. Thus Oscar can decrypt the message Alice sent, even though the cryptosystem may be "secure".
5. Illustrate the attack by computing $x$ by this method if $n=18721$, $b_{1}=43, b_{2}=7717, y_{1}=12677$ and $y_{2}=14702$.

A6.
We recall the three equations from the problem description:

$$
\begin{align*}
& c_{1}=b_{1}^{-1} \bmod b_{2}  \tag{4}\\
& c_{2}=\left(c_{1} b_{1}-1\right) / b_{2}  \tag{5}\\
& x_{1}=y_{1}^{c_{1}}\left(y_{2}^{c_{2}}\right)^{-1} \bmod n \tag{6}
\end{align*}
$$

1. In (4) both $b_{1}$ and $b_{2}$ are public and relatively prime thus Oscar can compute $c_{1}$.
2. In (5) note $c_{1} b_{1}=1+k b_{2}$ and thus $c_{1} b_{1}-1$ is divisible by $b_{2}$.
3. Rearranging (5) as $b_{1} c_{1}-b_{2} c_{2}=1$ and combining with (6) we get

$$
x_{1}=y_{1}^{c_{1}}\left(y_{2}^{c_{2}}\right)^{-1}=x^{b_{1} c_{1}}\left(x^{b_{2} c_{2}}\right)^{-1}=x^{b_{1} c_{1}-b_{2} c_{2}}=x
$$

4. $x_{1}$ is indeed the original plaintext $x$, which Oscar has recovered without knowledge of the private keys or factoring the modulus.

A6-b). We calculate

$$
\begin{aligned}
c_{1} & =43^{-1} \bmod 7717=2692 \\
c_{2} & =(2692 \cdot 43-1) / 7717=15 \\
x_{1} & =12677^{2692} \cdot\left(14702^{15}\right)^{-1} \bmod 18721 \\
& =13145 \cdot(3947)^{-1} \bmod 18721 \\
& =13145 \cdot 5668=15001 \bmod 18721
\end{aligned}
$$

and the plaintext $x_{1}=x=15001$. We can verify this as

$$
\begin{aligned}
15001^{43} \bmod 18721 & =12677=y_{1} \\
15001^{7717} \bmod 18721 & =14702=y_{2}
\end{aligned}
$$

