

- For each of the following 5-bit sequences determine its linear complexity and find one of the shortest LFSR that generates the sequence.

a) 0 0 1 1 1

b) 0 0 0 1 1

c) 1 1 1 0 0

- Find the shortest LFSR which generates all three sequences of problem 1.

- Let  $S$  be a sequence of bits with linear complexity  $L$ . Its complemented sequence  $\bar{S}$  is the sequence obtained from  $S$  by complementing its bits, that is, by adding 1 *modulo* 2 to each bit.

a) Show that  $LC(\bar{S}) \leq L + 1$ .

b) Show that  $LC(\bar{S}) = L - 1$ , or  $L$ , or  $L + 1$ .

- Use the Berlekamp-Massey Algorithm to find the shortest LFSR that generates the sequence:

0 0 1 0 1 0 1 1 1 1 1 0 0 .

Is this LFSR uniquely determined?

- Consider the 4-bit to 4-bit permutation  $\pi_S$  defined as follows:

0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
3	F	0	6	A	1	D	8	9	4	5	B	C	7	2	E

(This is the fourth row of the DES S-box  $S_4$ .) Denote by  $(x_1, x_2, x_3, x_4)$  and by  $(y_1, y_2, y_3, y_4)$  the input bits and output bits respectively. Find the output bit  $y_j$  for which the bias of  $x_1 \oplus x_2 \oplus x_3 \oplus x_4 \oplus y_j$  is the largest.

- Suppose that  $\mathbf{X}_1$  and  $\mathbf{X}_2$  are independent random variables defined on the set  $\{0, 1\}$ . Let  $\epsilon_i$  denote the bias of  $\mathbf{X}_i$ ,  $\epsilon_i = \Pr[\mathbf{X}_i = 0] - \frac{1}{2}$ , for  $i = 1, 2$ . Prove that if the random variables  $\mathbf{X}_1$  and  $\mathbf{X}_1 \oplus \mathbf{X}_2$  are independent, then  $\epsilon_2 = 0$  or  $\epsilon_1 = \pm \frac{1}{2}$ . (Hint: If the random variables  $\mathbf{X}_1$  and  $\mathbf{X}_1 \oplus \mathbf{X}_2$  are independent, then Piling-up lemma can be used to compute the bias of the  $\oplus$ -sum of these random variables.)