## T-79.503 Foundations of Cryptology

Homework 5
October 15, 2003

1. Given a positive integer $r$ and a combiner function $f: \mathbb{Z}_{26} \times \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}$ we define a kind of Feistel cipher as follows:

$$
\begin{aligned}
L_{i} & =R_{i-1}, \\
R_{i} & =\left(L_{i-1}+f\left(R_{i-1}, K_{i}\right)\right) \bmod 26,
\end{aligned}
$$

where $K_{i} \in \mathbb{Z}_{26}$, and $i=1,2, \ldots, r$, and $L_{j}, R_{j} \in \mathbb{Z}_{26}, j=0,1,2, \ldots, r$. The plaintext is $\left(L_{0}, R_{0}\right)$ and the ciphertext is $\left(L_{r}, R_{r}\right)$.
Consider a case where $r=3$ and the combiner function $f$ is defined as $f(X, K)=$ $(X \times K) \bmod 26$. The plaintext is $(2,11)$ and the ciphertext is $(8,1)$. Apply the meet-in-the-middle solution to find the keys $K_{1}$ and $K_{3}$. (Create tables as depicted in the figure, and find $K_{1}$ and $K_{3}$ such that $D\left(K_{1}\right)=D\left(K_{3}\right)$.


Figure 1: Meet-in-the-Middle solution
2. (cf. Stinson Exercise 3.3) Consider a Feistel cipher, where the $i$ th round is defined as follows:

$$
\begin{aligned}
L_{i} & =R_{i-1} \\
R_{i} & =L_{i-1} \oplus F_{i}\left(R_{i-1} \oplus K_{i}\right),
\end{aligned}
$$

where $K_{i}$ is the round key and $F_{i}$ is the round function. Given a bit sequence $A$ we denote by $c(A)$ the bit sequence obtained by complementing the bits of $A$, for example, if $A=001$, then $c(A)=110$. Let $Y=\left(L_{r}, R_{r}\right)$ be the ciphertext obtained by encrypting the plaintext $X=\left(L_{0}, R_{0}\right)$ (= concatenation of $L_{0}$ and $\left.R_{0}\right)$ using the $r$-round Feistel cipher with round keys $K_{1}, K_{2}, \ldots, K_{r}$. Show that then the plaintext $c(X)$ encrypted using the round keys $c\left(K_{1}\right), c\left(K_{2}\right), \ldots, c\left(K_{r}\right)$ gives the ciphertext $c(Y)$.
3. (Stinson 3.11 a)) The DES S-box $S_{4}$ has some unusual properties:

| 7 | 13 | 14 | 3 | 0 | 6 | 9 | 10 | 1 | 2 | 8 | 5 | 11 | 12 | 4 | 15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 13 | 8 | 11 | 5 | 6 | 15 | 0 | 3 | 4 | 7 | 2 | 12 | 1 | 10 | 14 | 9 |
| 10 | 6 | 9 | 0 | 12 | 11 | 7 | 13 | 15 | 1 | 3 | 14 | 5 | 2 | 8 | 4 |
| 3 | 15 | 0 | 6 | 10 | 1 | 13 | 8 | 9 | 4 | 5 | 11 | 12 | 7 | 2 | 14 |

Prove that the second row of $S_{4}$ can be obtained from the first row by means of the following mapping:

$$
\left(y_{1}, y_{2}, y_{3}, y_{4}\right) \mapsto\left(y_{2}, y_{1}, y_{4}, y_{3}\right) \oplus(0,1,1,0)
$$

4. Consider the $S_{4}$. Let us set $a=101110$. Which values the difference $S_{4}(x \oplus a) \oplus S_{4}(x)$ takes as $x$ varies through the sixteen values $x=\left(1, x_{2}, x_{3}, x_{4}, x_{5}, 1\right)$ ?
5. (Stinson 3.7) Suppose a sequence of plaintext blocks, $x_{1}, x_{2}, \ldots, x_{n}$ yields the ciphertext sequence $y_{1}, y_{2}, \ldots, y_{n}$. Suppose that one ciphertext block, say $y_{i}$, is transmitted incorrectly (i.e., some 1's are changed to 0 's and vice versa). Show that the number of plaintex blocks that will be decrypted incorrectly is equal to 1 if ECB or OFB modes are used for encryption; and equal to two if CBC or CFB modes are used.
