## T-79.4301 Spring 2006 Parallel and Distributed Systems Tutorial 8 - Wed Apr 12, 2006 11:15 and Fri Apr 21, 2006 14:15

- 1. Find Kripke models  $M_a$ ,  $M_b$ ,  $M_c$  and  $M_d$  (with  $AP = \{p, q\}$ ) such that
  - a)  $M_a \models \mathbf{G} p$  and  $M_a \models \mathbf{G} (p \Rightarrow q)$
  - b)  $M_b \not\models \mathbf{G} p$  and  $M_b \models \mathbf{G} (p \lor \mathbf{Y} q)$
  - c)  $M_c \models \mathbf{G} (p \Rightarrow (q \mathbf{S} \neg p))$  and  $M_c \models \mathbf{G} (p \Rightarrow \mathbf{Y} \mathbf{Y} \neg p)$
  - d)  $M_d \not\models \mathbf{G} (p \mathbf{S} q)$  and  $M_d \models \mathbf{G} \mathbf{O} q$

(For two formulas  $\psi_1, \psi_2$ , a finite word  $\pi = x_0 x_1 x_2 \dots x_n \in (2^{AP})^*$ , and an index  $0 \le i \le n$ ,  $\pi^i \models \psi_1 \Rightarrow \psi_2$  holds iff  $\pi^i \models (\neg \psi_1) \lor \psi_2$ .)

- 2. Let  $\varphi = \mathbf{G}(alarm \Rightarrow \mathbf{O}(crash))$  be a past safety formula over the atomic propositions  $AP = \{alarm, crash\}$ . Give a deterministic finite state automaton  $\overline{\mathcal{S}}$  (over the alphabet  $2^{AP}$ ) which accepts the finite words that are counterexamples to the formula  $\varphi$ .
- 3. Demo exercise: Model the automaton  $\overline{S}$  in Promela by (mis-)using the never claim construction to observe the global bool variables alarm and crash, and to execute the Promela statement assert(false) when the automaton  $\overline{S}$  would accept the sequence observed so far.