

Please note the following: your answers will be graded only if you have passed all the three home assignments before the exam!

Assignment 1 (10p)

- (a) Define the following concepts: *formation tree*, *free variable occurrence*, and *unique names assumption*. ($3 \times 2p$)
- (b) What is meant by the notation $\phi \equiv \psi$?
Prove in detail that if $\models \phi \rightarrow \psi$ and $\models \neg\phi \rightarrow \neg\psi$, then $\phi \equiv \psi$.

Assignment 2 (10p) Prove the following claims using semantic tableaux:

- (a) $\not\models (\neg A \leftrightarrow B \vee C) \leftrightarrow (A \leftrightarrow \neg B \vee \neg C)$
- (b) $\models \exists x \forall y Q(x, y) \rightarrow \forall y \exists x Q(x, y)$

Tableau proofs must contain all intermediary steps !!!

Assignment 3 (10p) Derive a Prenex normal form and a clausal form (i.e. a set of clauses S) for the sentence

$$\neg \forall x \exists y (\exists z R(x, z) \rightarrow \exists v R(y, v)).$$

Try to make S as simple as possible. Prove that S is unsatisfiable using resolution.

Assignment 4 (10p) Let us represent any finite string consisting of letters a , b , and c using unary function symbols a , b , and c and a constant symbol e denoting the empty string. Thus, for instance, the term $a(x)$ denotes a string that starts with an a followed by a string x , and the ground term $b(a(b(a(e))))$ represents “*baba*”.

- (a) Define the predicate $L(x, y) =$ “string x strictly precedes string y in the lexicographic order” so that your definition covers all finite strings represented as described above.
- (b) Give a model $\mathcal{S} \models \Sigma$ for your definition Σ according to which it holds that

$$\Sigma \not\models \exists x \exists y (L(x, y) \wedge L(y, x)).$$

Assignment 5 (10p)

Explain how the *weakest precondition* B_1 of an if-statement

$$\text{if}(B) \text{ then } \{C_1\} \text{ else } \{C_2\}$$

can be formed given a postcondition B_2 for it.

Consider the following program Divide:

$$v = 0 ; z = x ; \text{while}(z \geq y) \{ z = z - y ; v = v + 1 \}.$$

Use weakest preconditions and a suitable invariant to establish

$$\models_p [\text{true}] \text{Divide} [v == x / y],$$

where x / y denotes the integer quotient when x is divided by y .