## Autumn 2004

## T-79.148 Introduction to Theoretical Computer Science Tutorial 2, 28–29 September Problems

Remember to enroll for the course using the TOPI registration system by 1 October. Registration is compulsory.

## Homework problems:

- 1. Let  $\Sigma = \{a, b\}$ . Give some examples of strings from each of the following languages (at least three strings per language):
  - (a)  $\{w \in \Sigma^* \mid w \text{ the number of } a$ 's in w is even and the number of b's is divisible of three};
  - (b)  $\{a^{2n}b^{3m} \mid n, m \ge 0\};$
  - (c)  $\{uvu^Rv^R \mid u, v \in \Sigma^*\};^1$
  - (d)  $\{w \in \Sigma^* \mid \exists u, v \in \Sigma^* \text{ s.t. } w = uu = vvv\}.$
- 2. The reversal of a string  $w \in \Sigma^*$ , denoted  $w^R$ , is defined inductively by the rules:
  - (i)  $\varepsilon^R = \varepsilon;$
  - (ii) if w = ua, where  $u \in \Sigma^*$  and  $a \in \Sigma$ , then  $w^R = au^R$ .

It was proved in class (cf. also Lewis & Papadimitriou, p. 43) that for any strings  $u, v \in \Sigma^*$  it is the case that  $(uv)^R = v^R u^R$ . Prove in a similar manner, by induction based on the above definition of reversal, the following facts:

(a)  $(w^R)^R = w;$ 

(b) 
$$(w^k)^R = (w^R)^k$$
, for any  $k \ge 0$ .

3. Prove that for any two countable sets  $A_1$  and  $A_2$ , their union  $A_1 \cup A_2$  is also countable. Deduce from this by induction that the same holds for the union of n countable sets  $A_1 \cup A_2 \cup \ldots \cup A_n$ , for any  $n \ge 2$ . (*Extra question:* Does the claim still hold if the number of sets to be combined is countably infinite, i.e. in the case  $A = A_1 \cup A_2 \cup \ldots$ , where each  $A_i$  is countable?)

## **Demonstration problems:**

- 4. Show that any alphabet  $\Sigma$  with at least two symbols is comparable to the binary alphabet  $\Gamma = \{0, 1\}$ , in the sense that strings over  $\Sigma$  can be easily encoded into strings over  $\Gamma$  and vice versa. How much can the length of a string change in your encoding? (I.e., if the length of a string  $w \in \Sigma^*$  is |w| = n symbols, what is the length of the corresponding string  $w' \in \Gamma^*$ ?) Could you design a similar encoding if the target alphabet consisted of only *one* symbol, e.g.  $\Gamma = \{1\}$ ?
- 5. Prove that the Cartesian product  $\mathbb{N} \times \mathbb{N}$  is countably infinite. (*Hint:* Think of the pairs  $(m, n) \in \mathbb{N} \times \mathbb{N}$  as embedded in the Euclidean (x, y) plane  $\mathbb{R}^2$ . Enumerate the pairs by diagonals parallel to the line y = -x.) Conclude from this result that also the set  $\mathbb{Q}$  of rational numbers is countably infinite.
- 6. Let S be an arbitrary nonempty set.
  - (a) Give some injective (i.e. one-to-one) function  $f: S \to \mathcal{P}(S)$ .
  - (b) Prove that there cannot exist an injective function  $g : \mathcal{P}(S) \to S$ . (*Hint:* Assume that such a function g existed. Consider the set  $R = \{s \in S \mid s \notin g^{-1}(s)\}$ , and denote r = g(R). Is it then the case that  $r \in R$ ?)

Observe, as a consequence of item (b), that the power set  $\mathcal{P}(S)$  of any countably infinite set S is uncountable.

<sup>&</sup>lt;sup>1</sup>For a definition of the notation  $w^R$  see Problem 2.