Introduction to Theoretical Computer Science T/Y Tutorial 4, 14–15 February Problems

Homework problems:

- 1. Give regular expressions describing the following languages:
 - (a) $\{w \in \{a,b\}^* \mid w \text{ contains } abb \text{ as a substring}\};$
 - (b) $\{w \in \{a,b\}^* \mid w \text{ contains either } abb \text{ or } bba \text{ (or both) as a substring}\};$
 - (c) $\{w \in \{0,1\}^* \mid w \text{ contains exactly two 0's}\};$
 - (d) $\{w \in \{0,1\}^* \mid w \text{ contains at least two 0's}\};$
 - (e) $\{w \in \{0,1\}^* \mid w \text{ contains an even number (possibly zero) of 0's}\};$
 - (f) $\{w \in \{0,1\}^* \mid w \text{ begins and ends with different symbols}\};$
 - (g) $\{w \in \{0,1\}^* \mid |w| = 1 \pmod{3}\}.$
- 2. (a) Construct in a systematic way (as described in your textbook) a nondeterministic finite automaton corresponding to the regular expression $((\varepsilon \cup 0)1)^*011^*$.
 - (b) Make your automaton deterministic.
 - (c) Describe the language in part (a) in English as simply as you can.
- 3. Give regular expressions describing the following languages:
 - (a) $\{w \in \{a,b\}^* \mid w \text{ does not contain } abb \text{ as a substring}\};$
 - (b) $\{w \in \{0,1\}^* \mid w \text{ contains an even number of 0's and an odd number of 1's}\}.$

(*Hint:* Design first a finite automaton for each of the languages, cf. problems 2/1(c) & 2/3, and convert these automata then in a systematic manner, as described in your textbook, into the corresponding regular expressions.)

Demonstration problems:

- 4. Simplify the following regular expressions (i.e., design simpler expressions describing the same languages):
 - (a) $(\emptyset^* \cup a)(a^*)^*(b \cup a)b^*$
 - (b) $(a \cup b)^* \cup \emptyset \cup (a \cup b)b^*a^*$
 - (c) $a(b^* \cup a^*)(a^*b^*)^*$
- 5. Determine whether the regular expressions $r_1 = b^* a(a^*b^*)^*$ and $r_2 = (a \cup b)^* a(a \cup b)^*$ describe the same language, by constructing the minimal deterministic finite automata corresponding to them.
- 6. Prove that if L is a regular language, then so is $L' = \{xy \mid x \in L, y \notin L\}$.