

SOLUTIONS

1. It follows from the assumption  $H(\mathbf{K}|\mathbf{PC}) = 0$  that  $H(\mathbf{PKC}) = H(\mathbf{PC}) + H(\mathbf{K}|\mathbf{PC}) = H(\mathbf{PC})$ . For any cryptosystem,  $H(\mathbf{PKC}) = H(\mathbf{PK}) = H(\mathbf{P}) + H(\mathbf{K})$  holds, as the ciphertext is uniquely determined by the plaintext and the key, which are independent. From this we get

$$H(\mathbf{C}) + H(\mathbf{P}|\mathbf{C}) = H(\mathbf{PC}) = H(\mathbf{PK}) = H(\mathbf{P}) + H(\mathbf{K}).$$

It follows that the cryptosystem achieves perfect secrecy, that is,  $H(\mathbf{P}|\mathbf{C}) = H(\mathbf{P})$ , if and only if  $H(\mathbf{K}) = H(\mathbf{C})$ .

2. (6 pts) We compute  $11^{-1} \bmod 2008 = 1643$  and multiply the first equation with it to get  $x \equiv 913 \bmod 2008$ . We divide the second equation by three to get  $x \equiv 0 \bmod 669$ . Then we use the Chinese Remainder Theorem. For that purpose we compute

$$\begin{aligned} 2008^{-1} \bmod 2007 &= 1^{-1} \bmod 2007 = 1 \text{ and} \\ 669^{-1} \bmod 2008 &= (3^{-1}2007)^{-1} \bmod 2008 = 3(-1) \bmod 2008 = 2005 \end{aligned}$$

making use of the fact that  $2007 = -1 \bmod 2008$ . Then we get

$$x = 913 \cdot 2005 \cdot 669 + 0 \cdot 1 \cdot 2008 = 854313 \bmod (669 \cdot 2008),$$

where  $669 \cdot 2008 = 1343352$ . Then all three solutions modulo  $2007 \cdot 2008$  are  $854313 + i \cdot 1343352$ ,  $i = 0, 1, 2$ , that is,  $x = 854313, 2197665$  and  $3541017$ .

3. (a) We check that the irreducible polynomial  $f(x) = x^3 + x^2 + 1$  does not divide the polynomial  $g(x) = x^4 + x^3 + x^2 + 1$ . Then we can calculate  $(\text{lcm}(f(x), g(x))) = f(x)g(x) = x^7 + 1$ .
- (b) By Theorem 2 (Lecture 4) all sum sequences  $S_1 + S_2$  can be generated using polynomial  $x^7 + 1$ , which clearly has exponent 7. By Theorem 3 the period of the sum sequence divide 7. Hence 7 is the largest possible period.

4. We make the table

$x_1$	$x_2$	$x_3$	$y_1$	$y_2$	$x_1 \oplus x_2 \oplus y_1$	$x_1 \oplus x_2 \oplus y_2$
0	0	0	0	0	0	0
0	0	1	1	0	1	0
0	1	0	0	1	1	0
0	1	1	1	1	0	0
1	0	0	0	0	1	1
1	0	1	1	1	0	0
1	1	0	1	1	1	1
1	1	1	0	0	0	0

Then calculating the biases we get

$$\Pr[x_1 \oplus x_2 \oplus y_1 = 0] - \frac{1}{2} = \frac{4}{8} - \frac{1}{2} = 0 \text{ and}$$

$$\Pr[x_1 \oplus x_2 \oplus y_2 = 0] - \frac{1}{2} = \frac{6}{8} - \frac{1}{2} = \frac{1}{4} \neq 0.$$

The correct reply is:  $y_2$ .