

RSA Cryptosystem

$n = pq$ where p and q are two different large primes

$$\phi(n) = (p-1)(q-1)$$

a decryption exponent (private)

b encryption exponent (public)

$$ab \equiv 1 \pmod{\phi(n)}$$

RSA operation:

$$(m^b)^a \equiv m \pmod{n}$$

for all m , $0 \leq m < n$.

Wiener's result: It is insecure to select a shorter than about $\frac{1}{4}$ of the length of n .

RSA Equation

$$ab - k\phi(n) = 1$$

for some k where only b is known.

Additional information: $pq = n$ is known and $q < p < 2q$

$$n > \phi(n) = (p-1)(q-1) = pq - p - q + 1 \geq n - 3\sqrt{n}$$

Also we know that $a, b < \phi(n)$, hence $k < a$.

Wiener (1989) showed how to exploit this information to solve for a and all other parameters k, p and q , if a is sufficiently small.

Wiener's method is based on continued fractions.

Continued Fractions

Every rational number t has a unique representation as a finite chain of fractions

$$q_1 + \frac{1}{q_2 + \frac{1}{q_3 + \frac{1}{\ddots + \frac{1}{q_{m-1} + \frac{1}{q_m}}}}}$$

and we denote $t = [q_1 q_2 q_3 \dots q_{m-1} q_m]$. The rational number $t_j = [q_1 q_2 q_3 \dots q_j]$ is called the j^{th} convergent of t . For $t = u/v$, just run the Euclidean algorithm to find the q_i , $i = 1, 2, \dots, m$.

Fundamental Lemma

Theorem 5.14 *Suppose that $\gcd(u,v) = \gcd(c,d) = 1$ and*

$$\left| \frac{u}{v} - \frac{c}{d} \right| < \frac{1}{2d^2}.$$

Then c/d is one of the convergents of the continued fraction expansion of u/v .

Recall the RSA problem: $ab - k\phi(n) = 1$

Write it as:

$$\frac{b}{\phi(n)} - \frac{k}{a} = \frac{1}{a\phi(n)}$$

Then, if $2a < \phi(n)$, then k/a is a convergent of $b/\phi(n)$.

Wiener's Theorem

If in RSA cryptosystem

$$a < \frac{1}{3}\sqrt[4]{n},$$

that is, the length of the private exponent a is less than about one fourth of the length of n , then a can be computed in polynomial time with respect to the length of n .

Proof. First we show that k/a can be computed as a convergent of b/n , based on Euclidean algorithm, which is polynomial time. To see this, we estimate:

$$\left| \frac{b}{n} - \frac{k}{a} \right| = \left| \frac{ab - kn}{an} \right| = \left| \frac{1 + k\phi(n) - kn}{an} \right| \leq \frac{3k}{a\sqrt{n}} < \frac{3}{\sqrt{n}} < \frac{1}{2a^2}.$$

Wiener's Algorithm

Then the convergents $c_j/d_j = [q_1 q_2 q_3 \dots q_j]$ of b/n are computed. For the correct convergent $k/a = c_j/d_j$ we have

$$bd_j - c_j\phi(n) = 1.$$

For each convergent one computes

$$n' = (d_j b - 1) / c_j$$

and checks if $n' = \phi(n)$. Note that $p + q = n - \phi(n) + 1$.

Then if $n' = \phi(n)$, the equation

$$x^2 - (n - n' + 1)x + n = 0$$

has two positive integer solutions p and q .