

1. Show that the multiplicative order of  $g = 8 = 2^3$ , in modulo 19 arithmetic, is equal to 6.
2. Consider polynomial arithmetic with polynomial  $x^3 + x + 1$  on the set of three-bit integers. Determine the discrete logarithm of  $6 = 110$  to the base  $2 = 010$ .
3. Consider the polynomial arithmetic of four-bit numbers using the polynomial  $x^4 + x + 1$ . Then, all non-zero numbers form a cyclic group of order 15. Alice and Bob use Diffie-Hellman in this cyclic group with the generator element  $g = 3 = 0011$ . Alice's secret exponent  $a = 7$  and Bob's secret exponent  $b = 5$ . Compute the Diffie-Hellman key  $K$ .
4. Alice is using a toy version of the DSS signature scheme with a prime modulus  $p = 47$  and generator  $g = 2$  of order  $q = 23$ . By accident, Alice generates signatures for two different messages with the same per-messages random number  $k$ . The hash codes of the two signed messages are 2 and 3 and the signatures are  $(4, 21)$  and  $(4, 19)$ , respectively. Compute Alice's private key.
5. Alice and Bob use Diffie-Hellman Key Exchange to establish a shared secret key for their encrypted email application. After they computed the shared secret Diffie-Hellman key  $K = g^{ab}$  they verify that they have got the same value for  $K$ . For this purpose they compute a four digit hash value  $h(K)$  and compare the computed values by phone to ensure that they are equal. Suppose there is a man-in-the-middle  $C$ . Show how  $C$  can perform the man-in-the-middle attack in such a way that  $h(g^{ac}) = h(g^{bd})$ , with some values  $c$  and  $d$  chosen by  $C$ , that is, the attack is successful and remains undetected. What kind of computations  $C$  must do, and what is the expected amount of computations it must do to make the attack succeed? Consider the following two cases:
  - (a) Alice sends her public value first, and Bob sends his public value only after he received Alice's value.
  - (b) Alice and Bob send their public values in any order.