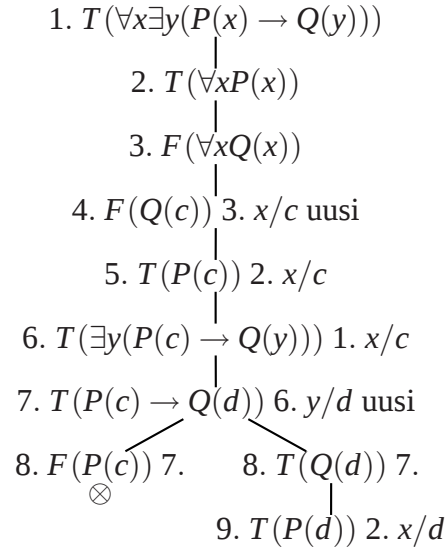


T-79.3001 Logic in Computer Science: Foundations **Spring 2009**
Exercise 9 ([Nerode and Shore, 1997], Predicate Logic, Chapters 4 and 9)
April 2 – 6, 2009

Solutions to demonstration problems

Solution to Problem 4

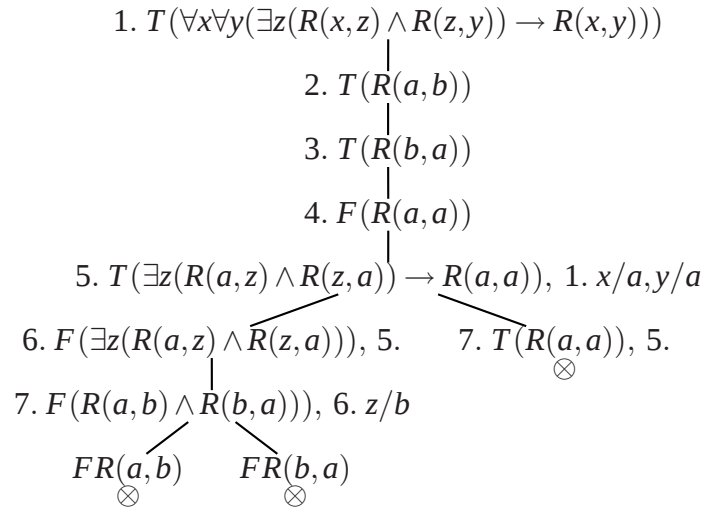
a) Tableau proof:



It seems that the tableau cannot be finished. We read a counter-example $\mathcal{S}$ from an open branch: domain $U = \{1, 2\}$, interpretations for constants $c^{\mathcal{S}} = 1$ and $d^{\mathcal{S}} = 2$, and interpretations for predicates $P^{\mathcal{S}} = \{1, 2\}$ and $Q^{\mathcal{S}} = \{2\}$.

Since the *tableau is not finished*, we need to check the counter-example. Now, we get $\mathcal{S} \models \forall x \exists y (P(x) \rightarrow Q(y))$, $\mathcal{S} \models \forall x P(x)$ and $\mathcal{S} \not\models \forall x Q(x)$ for $\mathcal{S}$.

b) Tableau proof:



All the branches in the tableau are contradictory and thus the claim holds.

Solution to Problem 5

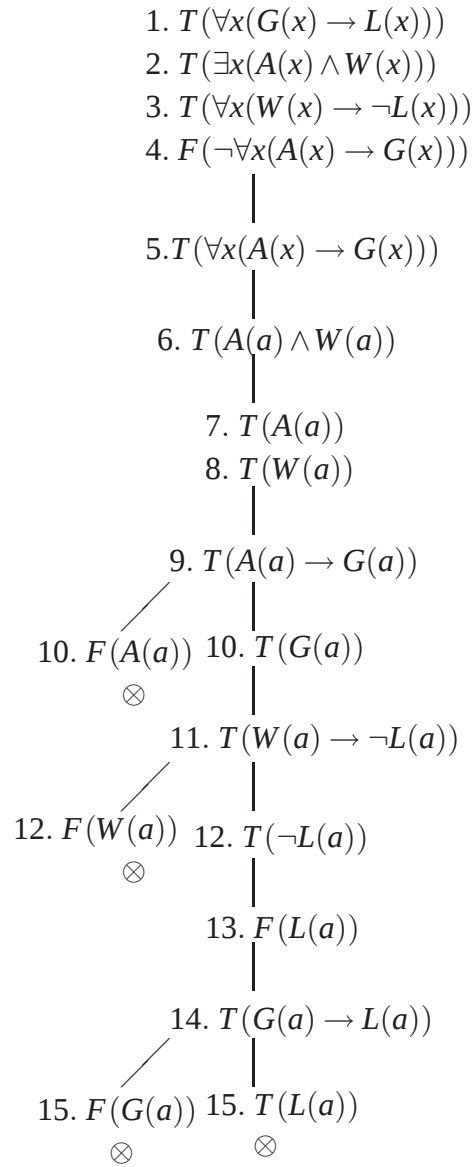
We choose the following predicates:

$$\begin{aligned} G(x) &= \text{“}x \text{ is guilty”}, \\ L(x) &= \text{“}x \text{ is liar”}, \\ A(x) &= \text{“}x \text{ is accused”}, \text{ and} \\ W(x) &= \text{“}x \text{ is witness”}. \end{aligned}$$

The sentences are:

- (i) $\forall x(G(x) \rightarrow L(x))$,
- (ii) $\exists x(A(x) \wedge W(x))$, ja
- (iii) $\forall x(W(x) \rightarrow \neg L(x))$.

and we want to show that $\neg \forall x(A(x) \rightarrow G(x))$. The tableaux proof is as follows.



Solution to Problem 6

We use the following predicates:

$$\begin{aligned} T(x,y) &= \text{“brick } x \text{ is on brick } y\text{”, and} \\ P(x) &= \text{“brick } x \text{ is on the table”}. \end{aligned}$$

The set of sentences is:

$$\{\forall x(\exists y T(x,y) \rightarrow \neg P(x)), \forall x(P(x) \vee \exists y T(x,y)), \\ \forall x \forall y(\exists z T(y,z) \rightarrow \neg T(x,y))\}$$

and we want to show that $\forall x \forall y(T(x,y) \rightarrow P(y))$.

Tableau proof:

$$\begin{array}{l} 1. T(\forall x(\exists y T(x,y) \rightarrow \neg P(x))) \\ \quad | \\ 2. T(\forall x(P(x) \vee \exists y T(x,y))) \\ \quad | \\ 3. T(\forall x \forall y(\exists z T(y,z) \rightarrow \neg T(x,y))) \\ \quad | \\ 4. F(\forall x \forall y(T(x,y) \rightarrow P(y))) \\ \quad | \\ 5. F(\forall y(T(c,y) \rightarrow P(y)))^{4. \text{ } x/c \text{ uusi}} \\ \quad | \\ 6. F(T(c,d) \rightarrow P(d))^{5. \text{ } y/d \text{ uusi}} \\ \quad | \\ 7. T(T(c,d))^{6.} \\ \quad | \\ 8. F(P(d))^{6.} \\ \quad | \\ 9. T(P(d) \vee \exists y T(d,y)) \\ \quad / \quad \backslash \\ 10. T(P(d))^{9.} \quad 10. T(\exists y T(d,y))^{9.} \\ \quad \otimes \quad \quad | \\ \quad \quad 11. T(T(d,e))^{10. \text{ } y/e \text{ uusi}} \\ \quad \quad | \\ \quad \quad 12. T(\exists z T(d,z) \rightarrow \neg T(c,d))^{3. \text{ } x/c, y/d} \\ \quad \quad / \quad \backslash \\ \quad \quad 13. F(\exists z T(d,z))^{12.} \quad 13. T(\neg T(c,d))^{12.} \\ \quad \quad | \quad \quad | \\ \quad \quad 14. F(T(d,e))^{13. \text{ } z/e} \quad 14. F(T(c,d))^{13.} \\ \quad \quad \otimes \quad \quad \otimes \end{array}$$

Note: 1) can be equivalently stated as $\forall x \forall y(T(x,y) \rightarrow \neg P(x))$ and 3) as $\forall x \forall y \forall z(T(y,z) \rightarrow \neg T(x,y))$. How would the tableau look if you used these sentences?