

Please note the following: your answers will be graded only if you have passed all the three home assignments before the exam!

Assignment 1 (10p)

- (a) Define the following concepts: *model*, *ground term*, and *unifier*. (3 × 2p)
- (b) What is meant by the notation $\phi \equiv \psi$?
Prove in detail that if $\phi \equiv \psi$ and $\psi \equiv \chi$, then $\phi \equiv \chi$. (4p)

Assignment 2 (10p) Prove the following claims using semantic tableaux:

- (a) $\models (A \rightarrow B) \wedge (B \rightarrow C) \wedge (C \rightarrow A) \rightarrow (\neg B \rightarrow \neg C)$
- (b) $\models \forall y \forall x (R(x, y) \rightarrow P(x)) \rightarrow \forall x (\exists y R(x, y) \rightarrow P(x))$

Tableau proofs must contain all intermediary steps !!!

Assignment 3 (10p) Derive a Prenex normal form and a clausal form (i.e. a set of clauses S) for the sentence

$$\neg(\forall x \exists y (P(x) \vee Q(y)) \rightarrow \exists y \forall x (P(x) \vee Q(y))).$$

Make S as simple as possible. Prove that S is unsatisfiable using resolution.

Assignment 4 (10p) Let us represent strings “”, “a”, “b”, “aa”, “ab”, “ba”, “bb”, ... that consist of letters a ja b using ground terms

$$e, a(e), b(e), a(a(e)), a(b(e)), b(a(e)), b(b(e)), \dots,$$

built of a constant symbol e , which represents the empty string “”, and unary functions $a(x)$ and $b(x)$, that append the respective letter a or b at the beginning of a string x . Thus $a(b(e))$ is interpreted as $a(b(\text{“”})) = a(\text{“b”}) = \text{“ab”}$.

- (a) Define predicate $P(x) = \text{“the letter } a \text{ occurs an even number of times in the string } x\text{”}$ using predicate logic so that your definition covers all finite strings represented as explained above.
- (b) Give a model $\mathcal{S} \models \Sigma$ of your definition Σ on the basis of which it holds that

$$\Sigma \not\models P(b(a(b(e))))).$$

Assignment 5 (10p)

Explain how the *weakest precondition* B_1 of an if-statement

$$\text{if}(B) \text{ then } \{C_1\} \text{ else } \{C_2\}$$

can be formed given a postcondition B_2 for it.

Consider the following program Minus:

$$v = x ; z = y ; \text{while}(! (z == 0)) \{ z = z - 1 ; v = v - 1 \}.$$

Use weakest preconditions and a suitable invariant to establish

$$\models_p [\text{true}] \text{Minus} [v == x - y].$$