

**Please note the following: your answers will be graded only if you have passed all the three home assignments before the exam!**

**Assignment 1** (10p)

- (a) Define the following concepts: *model*, *ground term*, and *unifier*. (3 × 2p)
- (b) What is meant by the notation  $\phi \equiv \psi$ ?  
Prove in detail that if  $\phi \equiv \psi$  and  $\psi \equiv \chi$ , then  $\phi \equiv \chi$ . (4p)

**Assignment 2** (10p) Prove the following claims using semantic tableaux:

- (a)  $\models (A \rightarrow B) \wedge (B \rightarrow C) \wedge (C \rightarrow A) \rightarrow (\neg B \rightarrow \neg C)$
- (b)  $\models \forall y \forall x (R(x, y) \rightarrow P(x)) \rightarrow \forall x (\exists y R(x, y) \rightarrow P(x))$

Tableau proofs must contain all intermediary steps !!!

**Assignment 3** (10p) Derive a Prenex normal form and a clausal form (i.e. a set of clauses  $S$ ) for the sentence

$$\neg(\forall x \exists y (P(x) \vee Q(y)) \rightarrow \exists y \forall x (P(x) \vee Q(y))).$$

Make  $S$  as simple as possible. Prove that  $S$  is unsatisfiable using resolution.

**Assignment 4** (10p) Let us represent strings “”, “a”, “b”, “aa”, “ab”, “ba”, “bb”, ... that consist of letters  $a$  ja  $b$  using ground terms

$$e, a(e), b(e), a(a(e)), a(b(e)), b(a(e)), b(b(e)), \dots,$$

built of a constant symbol  $e$ , which represents the empty string “”, and unary functions  $a(x)$  and  $b(x)$ , that append the respective letter  $a$  or  $b$  at the beginning of a string  $x$ . Thus  $a(b(e))$  is interpreted as  $a(b(\text{“”})) = a(\text{“b”}) = \text{“ab”}$ .

- (a) Define predicate  $P(x) = \text{“the letter } a \text{ occurs an even number of times in the string } x\text{”}$  using predicate logic so that your definition covers all finite strings represented as explained above.
- (b) Give a model  $\mathcal{S} \models \Sigma$  of your definition  $\Sigma$  on the basis of which it holds that

$$\Sigma \not\models P(b(a(b(e)))).$$

**Assignment 5** (10p)

Explain how the *weakest precondition*  $B_1$  of an if-statement

$$\text{if}(B) \text{ then } \{C_1\} \text{ else } \{C_2\}$$

can be formed given a postcondition  $B_2$  for it.

Consider the following program Minus:

$$v = x ; z = y ; \text{while}(! (z == 0)) \{ z = z - 1 ; v = v - 1 \}.$$

Use weakest preconditions and a suitable invariant to establish

$$\models_p [\text{true}] \text{Minus} [v == x - y].$$