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# Local Search Algorithms for Random Satisfiability

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# Outline

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- Background
- WalkSAT and related algorithms
- Results and conjectures on WalkSAT
- Experiments on WalkSAT
- Record-to-Record Travel and variants
- Experiments on RRT
- Focused Metropolis Search
- Experiments on FMS
- Dynamics of FMS
- Analysis?



# Background

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Denote  $N$  = number of variables,  $M$  = number of clauses,  
 $\alpha = M/N$ .

Satisfiability transition at  $\alpha_c \approx 4.267$  (Mitchell et al. 1992, . . . ,  
Braunstein et al. 2002).

Good experiences with local search methods in the satisfiable  
region  $\alpha < \alpha_c$ : e.g. GSAT (Selman et al. 1992), WalkSAT (Selman  
et al. 1996). Experiments 1996:  $N \leq 2000$  at  $\alpha \approx \alpha_c$ .



# GSAT

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Selman et al. 1992 ... 1996.

Denote by  $E = E_F(s)$  the number of unsatisfied clauses in formula  $F$  under truth assignment  $s$ .

GSAT( $F$ ):

```
s = initial truth assignment;
while flips < max_flips do
  if s satisfies F then output s & halt, else:
  - find a variable x whose flipping causes
    largest decrease in E (if no decrease is
    possible, then smallest increase);
  - flip x.
```



# NoisyGSAT

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GSAT augmented by a fraction  $p$  of random walk moves.

NoisyGSAT( $F, p$ ):

$s$  = initial truth assignment;

while flips < max\_flips do

if  $s$  satisfies  $F$  then output  $s$  & halt, else:

- with probability  $p$ , pick a variable  $x$  uniformly at random and flip it;
- with probability  $(1-p)$ , do basic GSAT move:
  - find a variable  $x$  whose flipping causes largest decrease in  $E$  (if no decrease is possible, then smallest increase);
  - flip  $x$ .



# WalkSAT

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NoisyGSAT *focused* on the unsatisfied clauses.

WalkSAT( $F, p$ ):

$s$  = initial truth assignment;

while flips < max\_flips do

if  $s$  satisfies  $F$  then output  $s$  & halt, else:

- pick a random unsatisfied clause  $C$  in  $F$ ;
- if some variables in  $C$  can be flipped without breaking any presently satisfied clauses, then pick one such variable  $x$  at random; else:
- with probability  $p$ , pick a variable  $x$  in  $C$  at random;
- with probability  $(1-p)$ , pick an  $x$  in  $C$  that breaks a minimal number of presently satisfied clauses;
- flip  $x$ .



# WalkSAT vs. NoisyGSAT

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The focusing seems to be important: in the (unsystematic) experiments in Selman et al. (1996), WalkSAT outperforms NoisyGSAT by several orders of magnitude.



# Recent results and conjectures

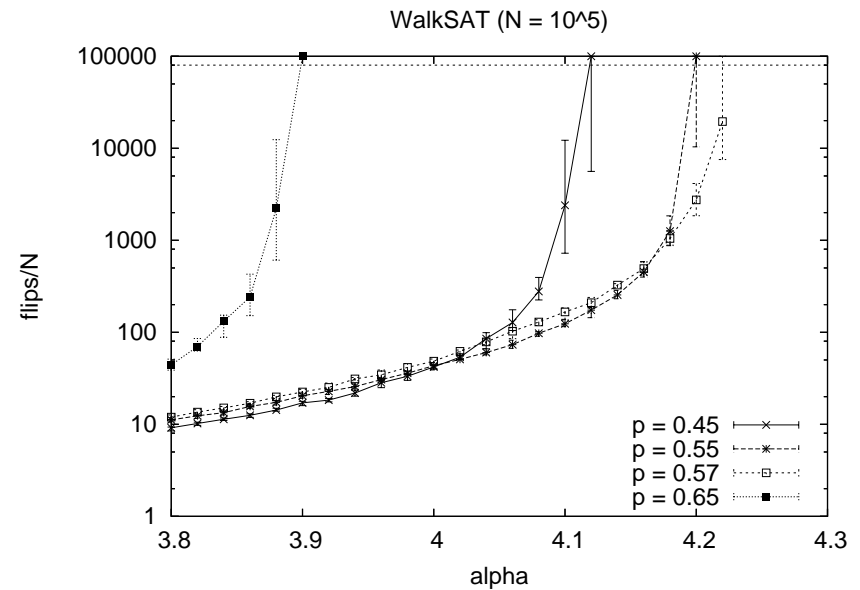
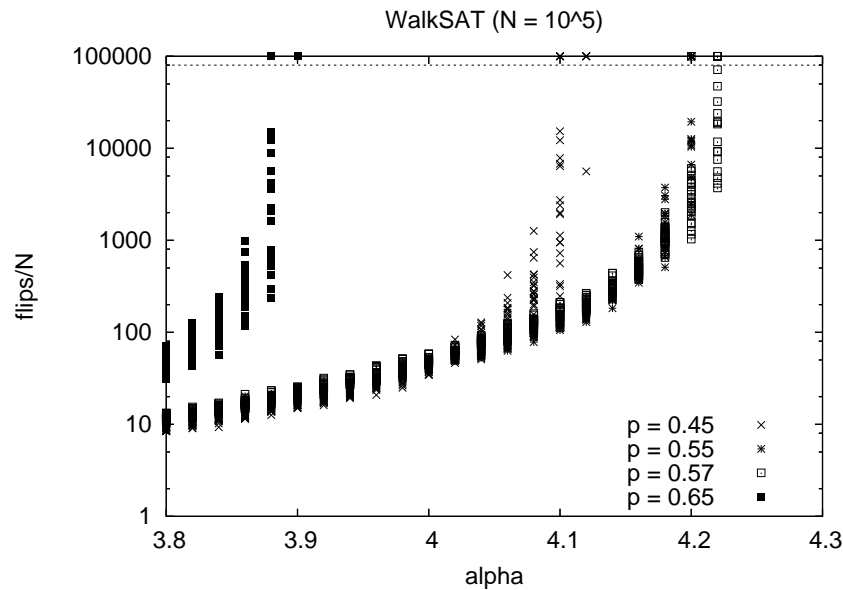
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- Barthele et al. (2003): numerical experiments with WalkSAT (mainly with  $p = 1$ , some also with  $p < 1$ ) at  $N = 50,000$ ,  $\alpha = 2.0 \dots 4.0$ . Observed transition in the dynamics at  $\alpha_{\text{dyn}} \approx 2.7 - 2.8$ . When  $\alpha < \alpha_{\text{dyn}}$ , satisfying assignments are found in linear time per variable (i.e. in a total of  $cN$  “flips”), when  $\alpha > \alpha_{\text{dyn}}$  exponential time is required.
- Similar results obtained by Semerjian & Monasson (2003), though with smaller experiments ( $N = 500$ ).
- Explanation: for  $\alpha > \alpha_{\text{dyn}}$  the search equilibrates at a nonzero energy level, and can only escape to a ground state through a large enough random fluctuation.
- Conjecture: no local search algorithm works in linear time beyond the clustering transition at  $\alpha_s \approx 3.92 - 3.93$  (Mézard, Monasson, Weigt et al.)





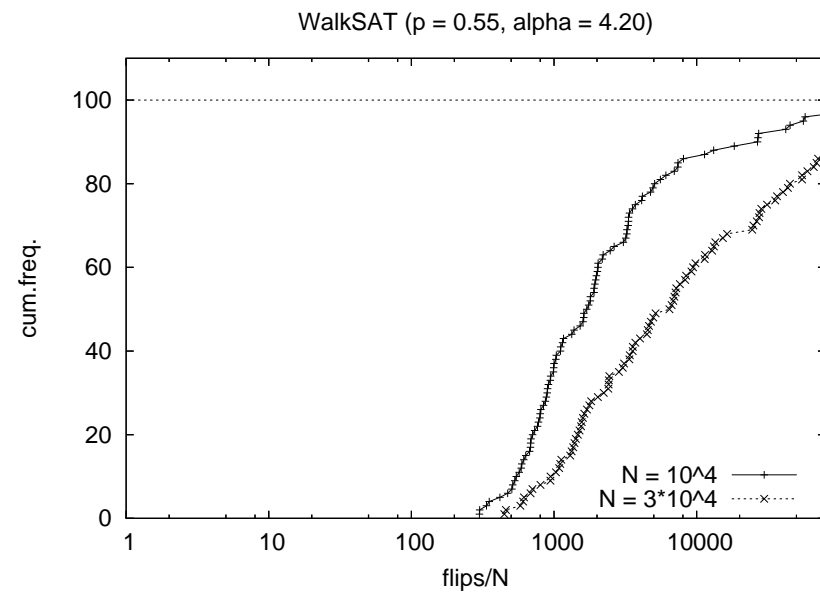
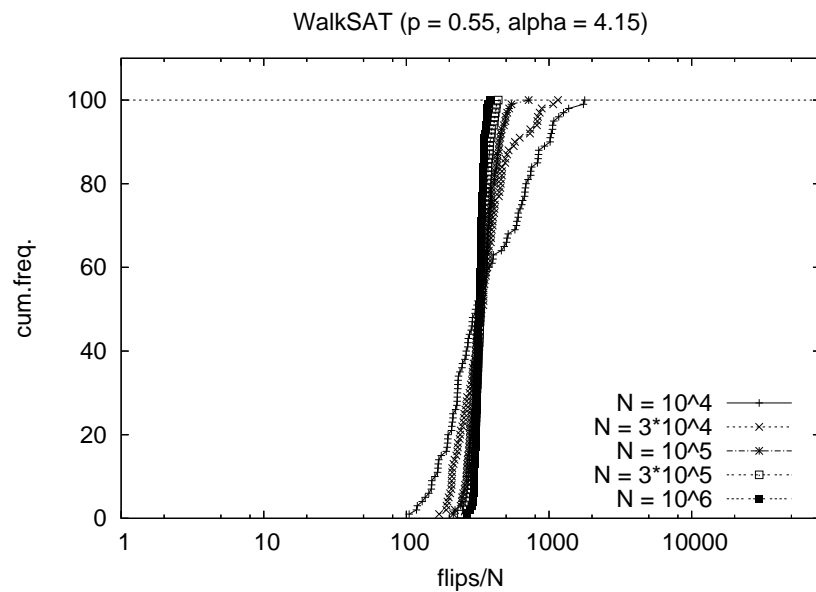
# WalkSAT experiments (3-SAT)



Normalised solution times for WalkSAT,  $\alpha = 3.8 \dots 4.3$ .  
Left: complete data; right: medians and quartiles.



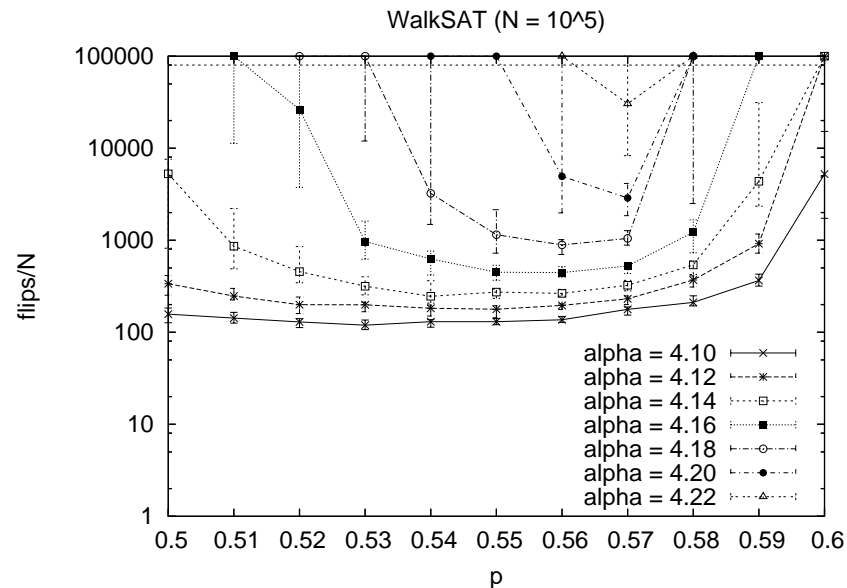
# WalkSAT linear scaling



Cumulative solution time distributions for WalkSAT with  $p = 0.55$ .



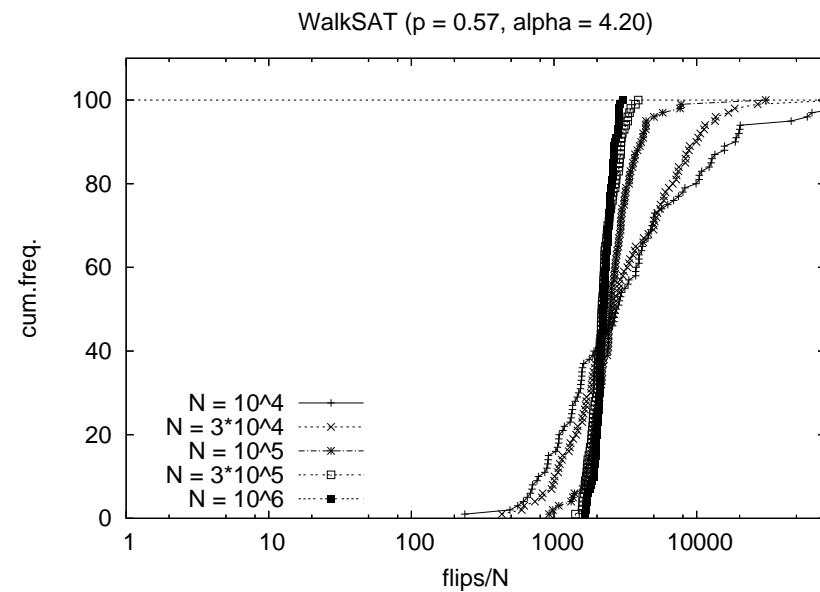
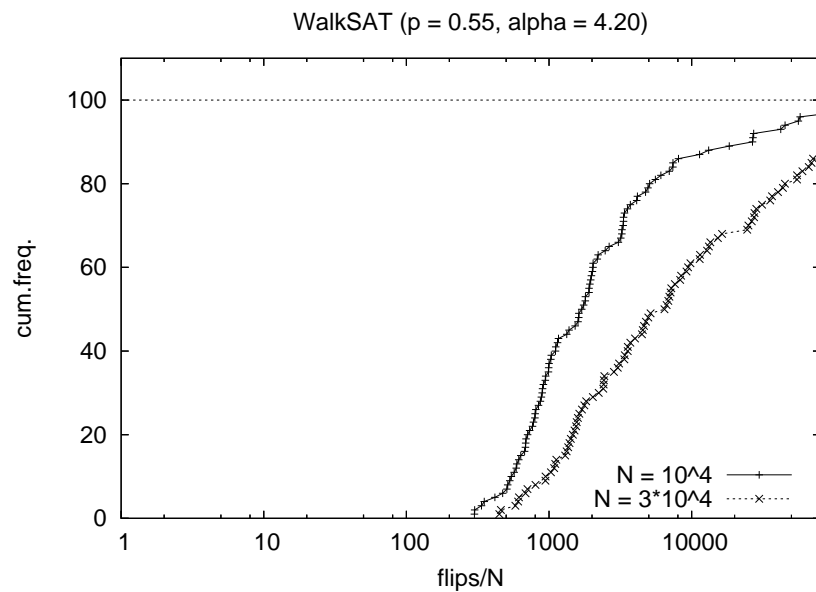
# WalkSAT optimal noise level?



Normalised solution times for WalkSAT with  $p = 0.50 \dots 0.60$ ,  
 $\alpha = 4.10 \dots 4.22$ .



# WalkSAT sensitivity to noise



Cumulative solution time distributions for WalkSAT at  $\alpha = 4.20$  with  $p = 0.55$  and  $p = 0.57$ .



# Record-to-Record Travel (RRT)

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Very simple stochastic local optimisation algorithm introduced by Dueck (1993). Dueck claimed good results on solving 442-city and 532-city TSP's; after that little used.

RRT( $E, d$ ):

```
s = initial feasible solution;
```

```
s* = s; E* = E(s);
```

```
while moves < max_moves do
```

```
  if s is a global min. of E then output s & halt,  
  else:
```

```
    pick a random neighbour s' of s;
```

```
    if E(s') <= E* + d then let s = s';
```

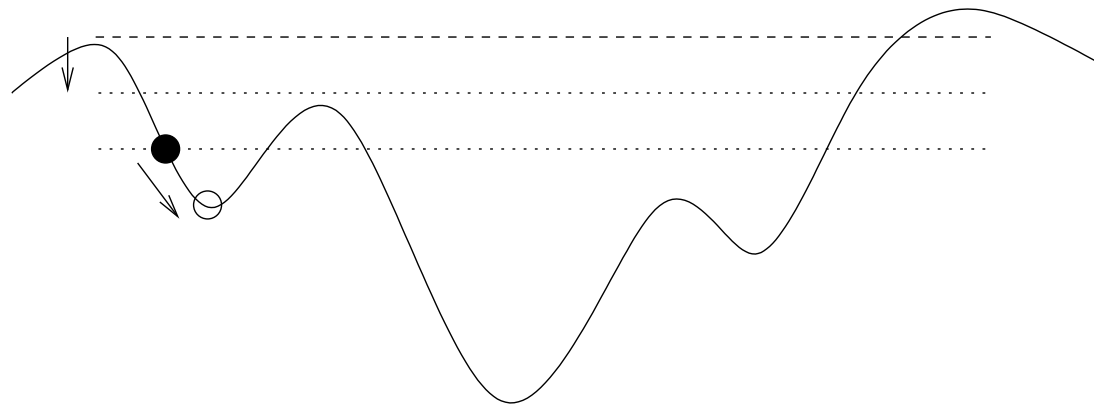
```
    if E(s') < E* then:
```

```
      s* = s'; E* = E(s').
```



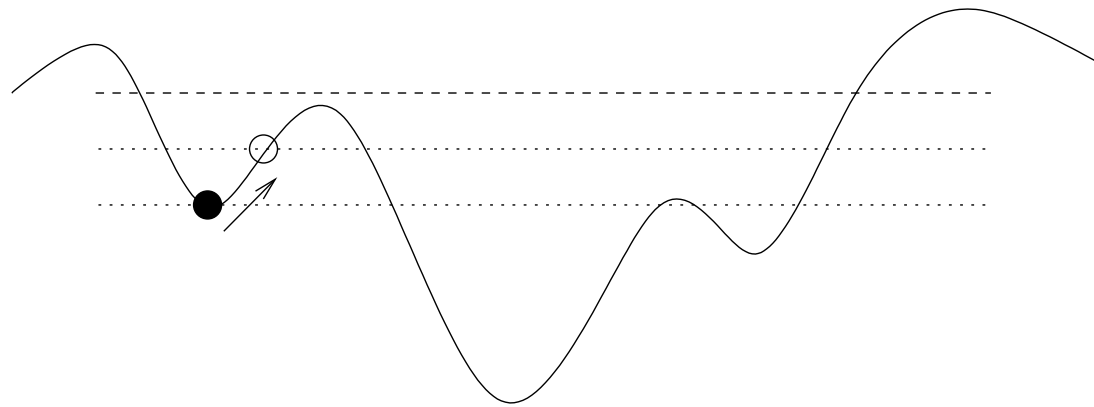
# RRT in action (d = 2)

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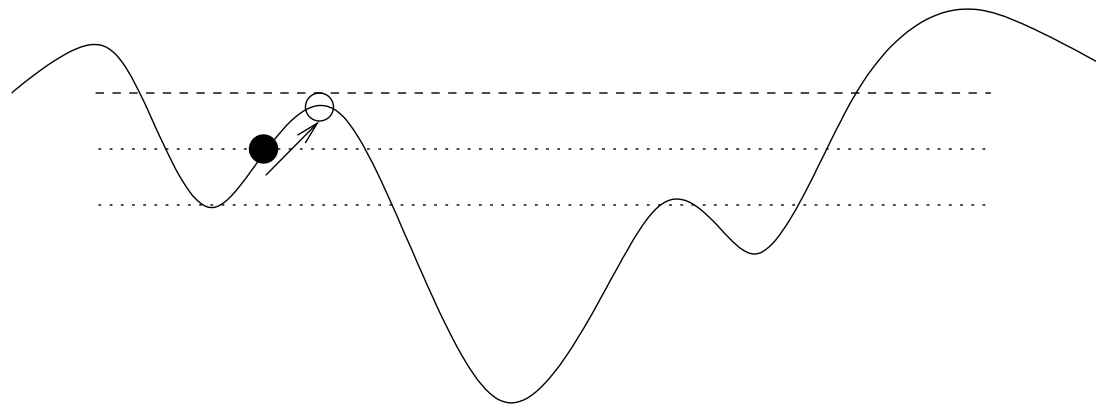
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# RRT in action (d = 2)

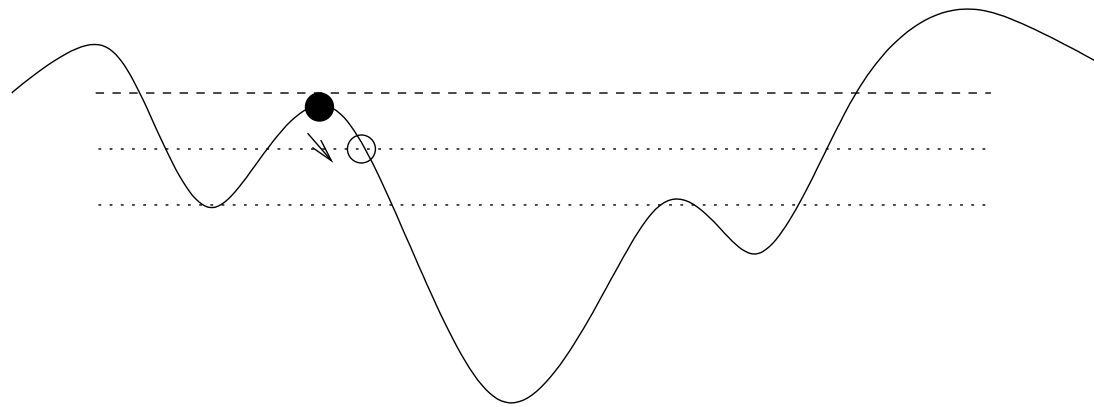
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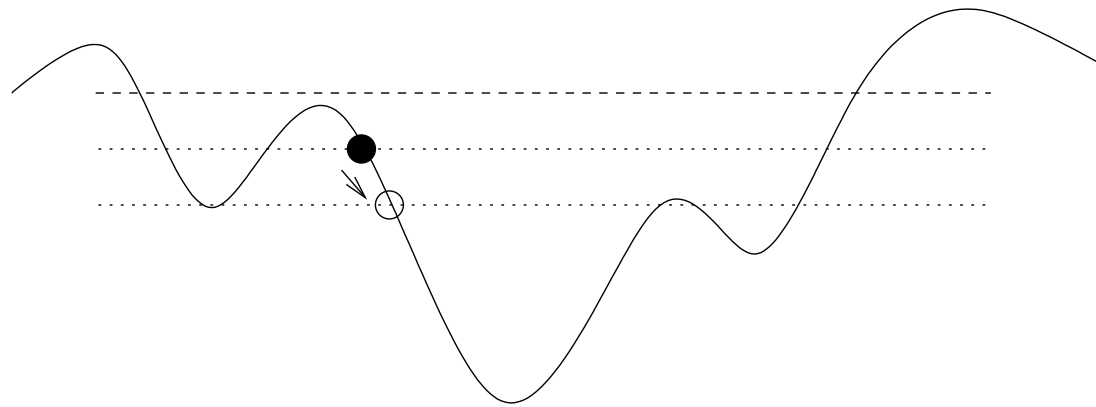
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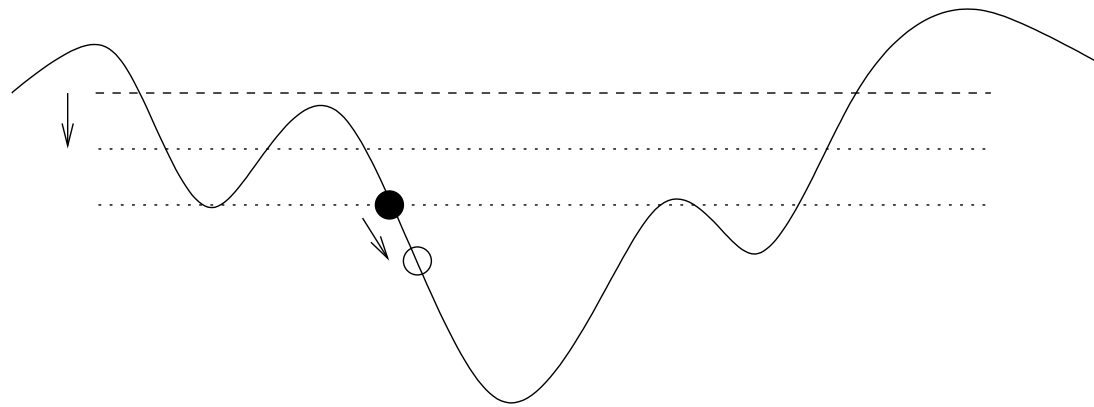
# RRT in action (d = 2)

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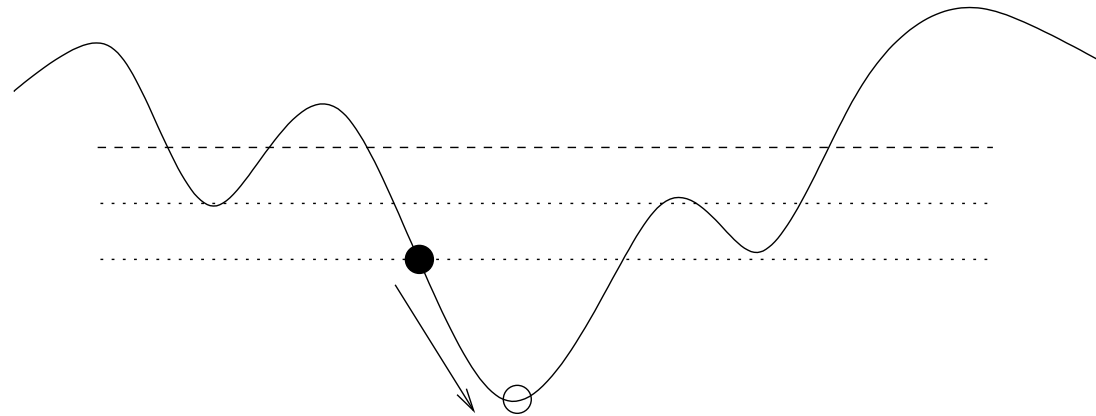
# RRT in action (d = 2)

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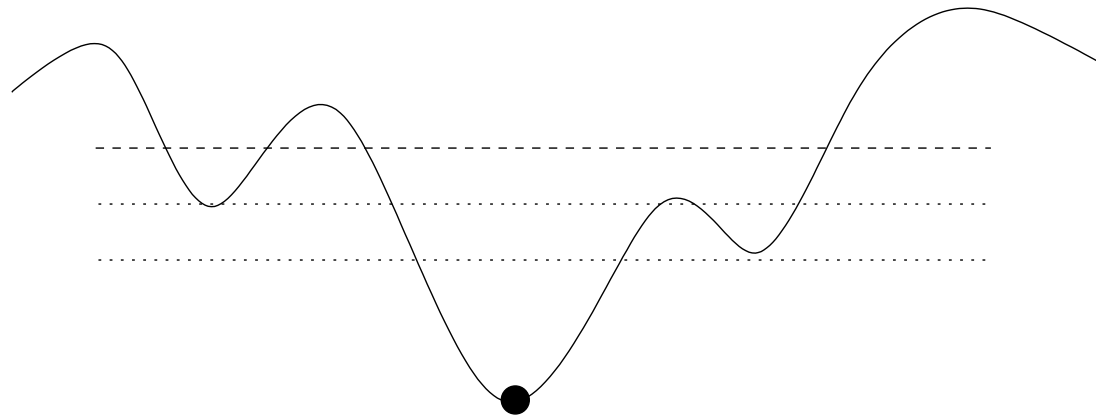
# RRT in action (d = 2)

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# RRT in action (d = 2)

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# Focused RRT

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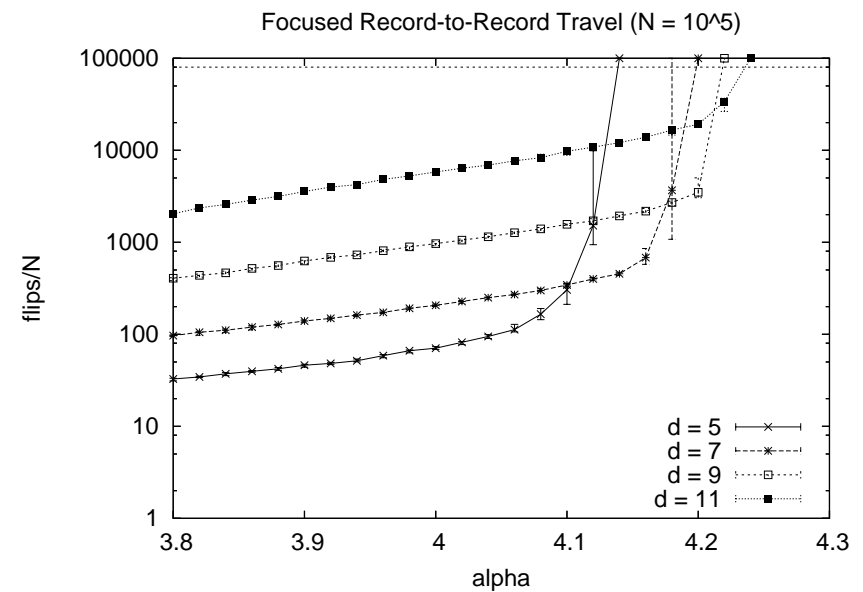
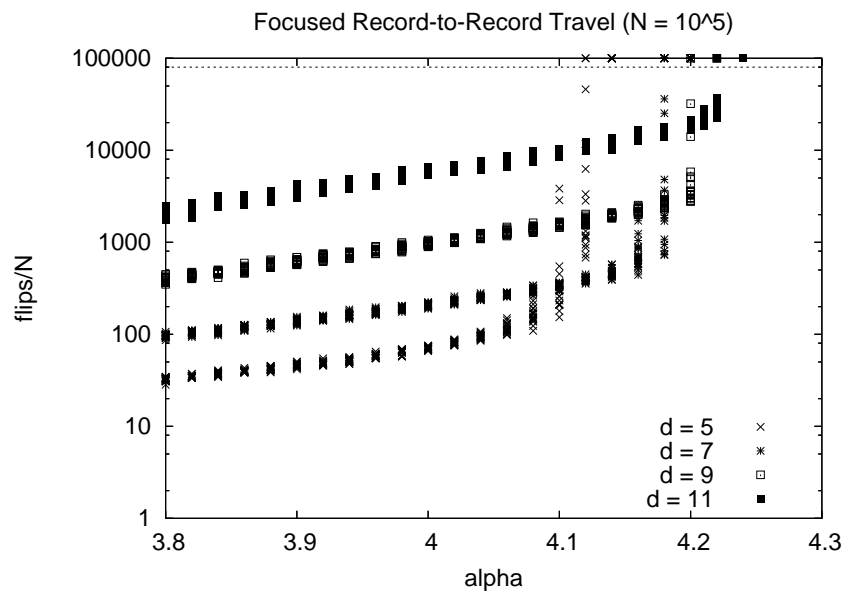
In applying RRT to SAT,  $E(s)$  = number of clauses unsatisfied by truth assignment  $s$ . Single-variable flip neighbourhoods.

*Focusing*: flipped variables chosen from unsatisfied clauses.  
(Precisely: one unsatisfied clause is chosen at random, and from there a variable at random.)

⇒ FRRT = focused RRT.



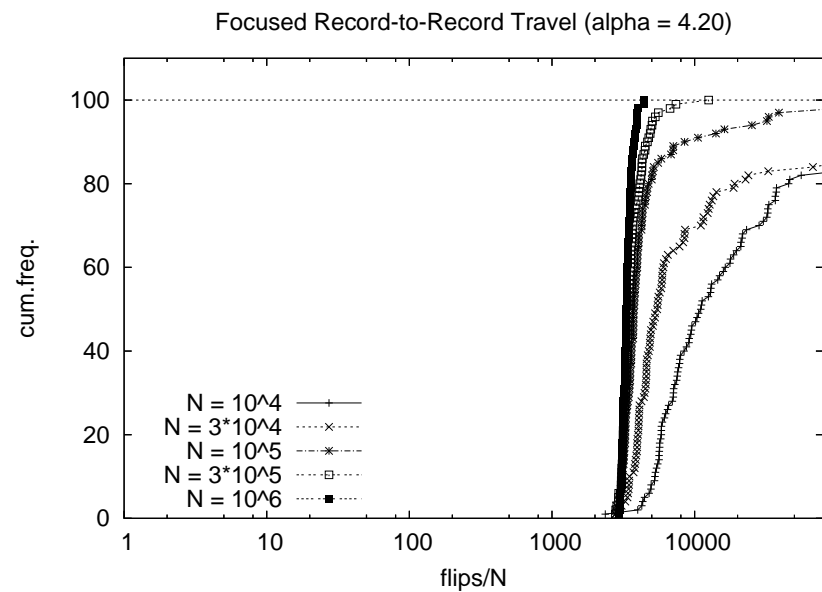
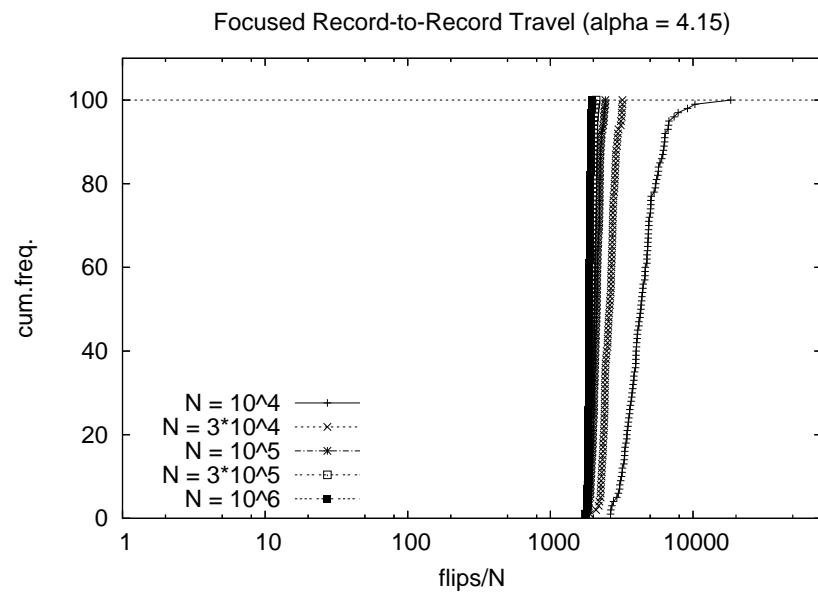
# FRRT experiments (3-SAT)



Normalised solution times for FRRT,  $\alpha = 3.8 \dots 4.3$ .  
Left: complete data; right: medians and quartiles.



# FRRT linear scaling

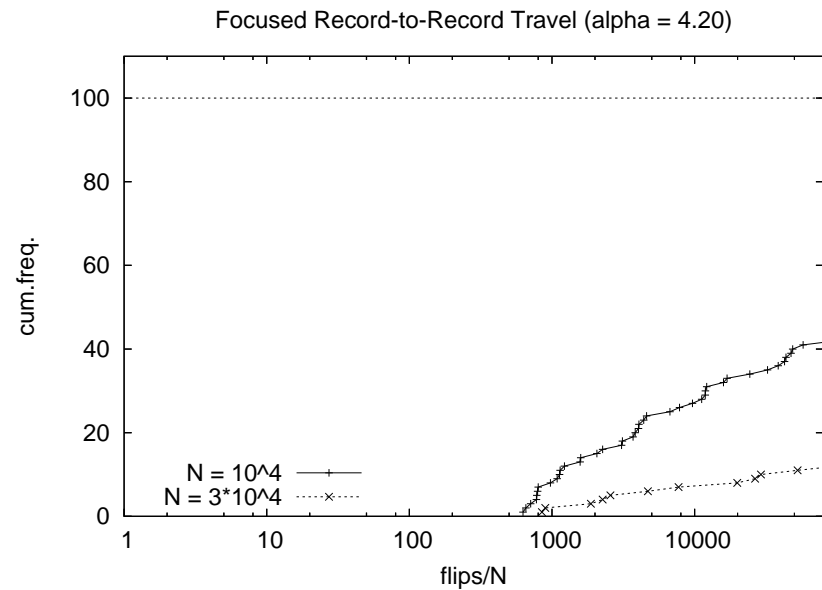
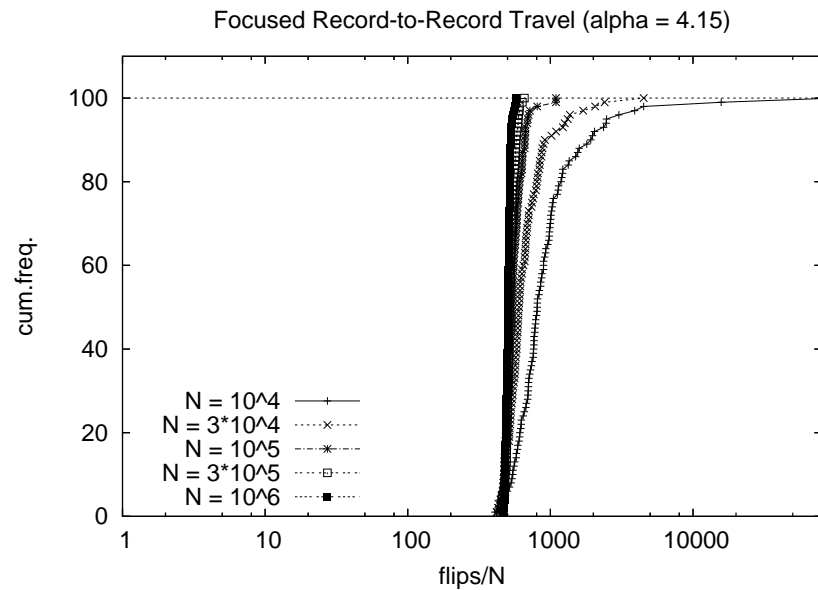


Cumulative solution time distributions for FRRT with  $d = 9$ .





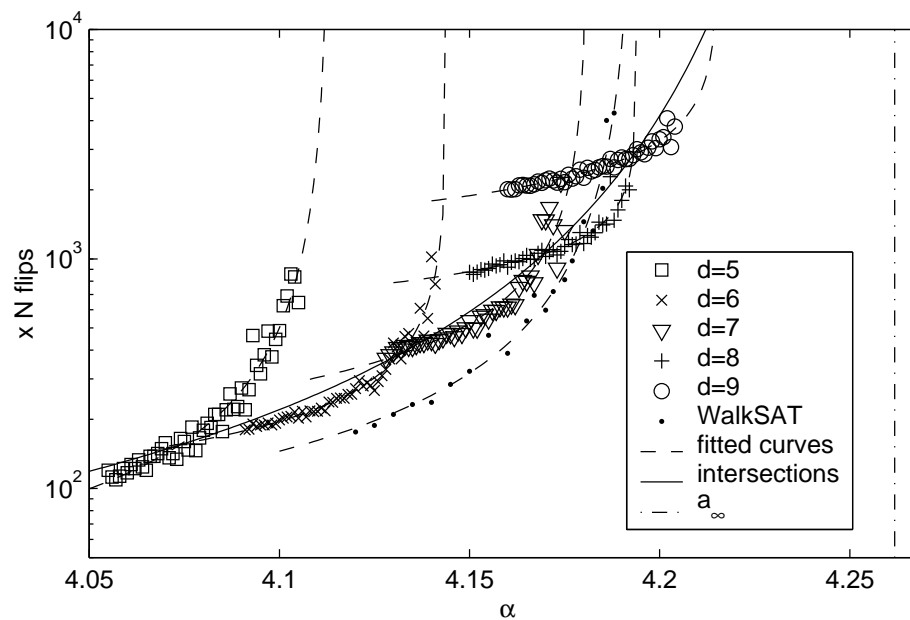
# FRRT linear scaling (cont'd)



Cumulative solution time distributions for FRRT with  $d = 7$ .



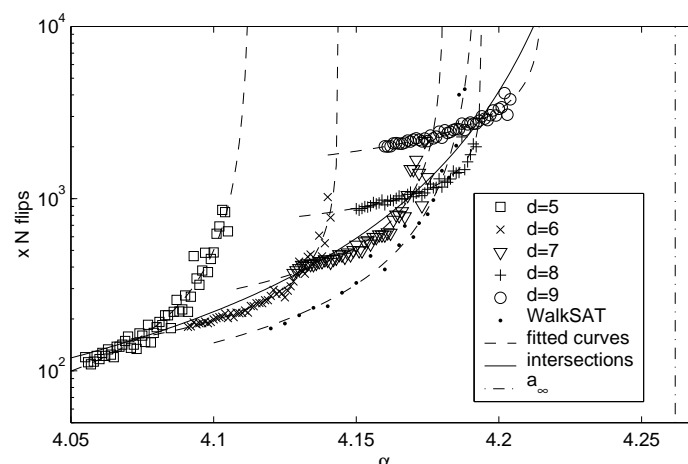
# FRRT qualitative observations



For each fixed  $d$  there seems to be a transition value  $\alpha_d$  s.th. for  $\alpha < \alpha_d$  the algorithm runs in linear time per variable, and for  $\alpha > \alpha_d$  requires exponential time per variable. (Empirical estimates:  $\alpha_5 \approx 4.11$ ,  $\alpha_6 \approx 4.14$ ,  $\alpha_7 \approx 4.18$ ,  $\alpha_8 \approx 4.19$ ,  $\alpha_9 \approx 4.21$ .)



# FRRT qualitative observations cont'd



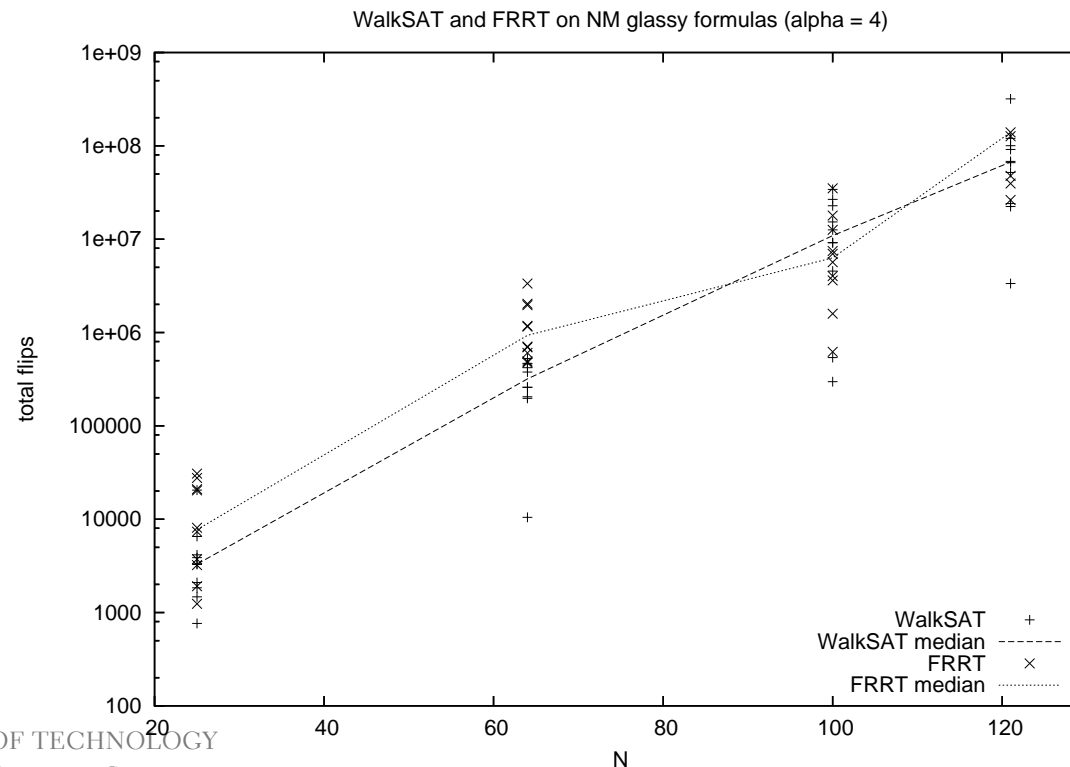
An empirical fit of the transition points  $\alpha_d$  suggests that for large  $d$  they converge towards  $\alpha_\infty \approx 4.26$ .

Comparative experiments using WalkSAT with near optimal parameter settings ( $p = 0.55$ ) yield estimate  $\alpha_{\text{dyn}} \approx 4.19$  for WalkSAT's transition point.



# WalkSAT & FRRT on structured problems

Jia, Moore & Selman (2004) tested WalkSAT and FRRT on highly structured “glassy” 3-SAT formulas. Here the number of variables is always of the form  $N = L \times L$ ; values of  $L = 5, 8, 10, 11, 16$  were tried out. At  $L = 16$  WalkSAT no longer converged; FRRT did, but only for  $d = 5$ .



# Focused Metropolis Search

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Arguably the most natural focused local search algorithm. Variable flip acceptance probabilities determined by a parameter  $\eta$ ,  $0 \leq \eta \leq 1$ .

FMS( $F, \eta$ ):

`s = initial truth assignment;`

`while flips < max_flips do`

`if s satisfies F then output s & halt, else:`

`pick a random unsatisfied clause C in F;`

`pick a variable x in C at random;`

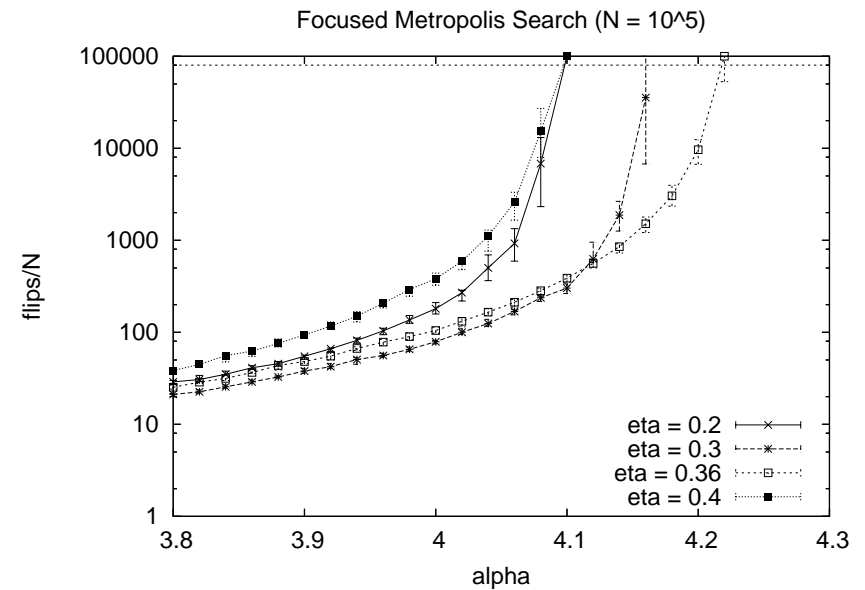
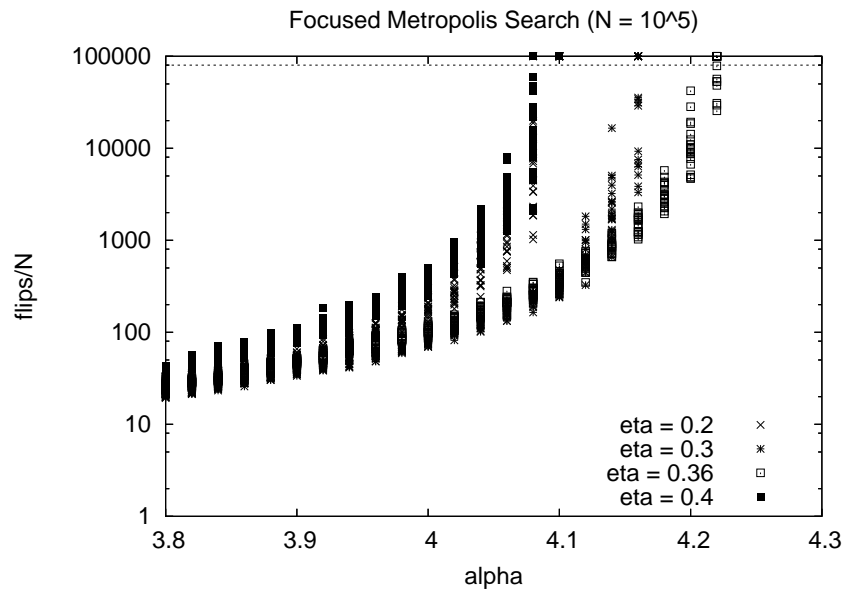
`let x' = flip(x), s' = s[x'/x];`

`if E(s') <= E(s) then flip x, else:`

`flip x with prob.  $\eta^{E(s')-E(s)}$ .`



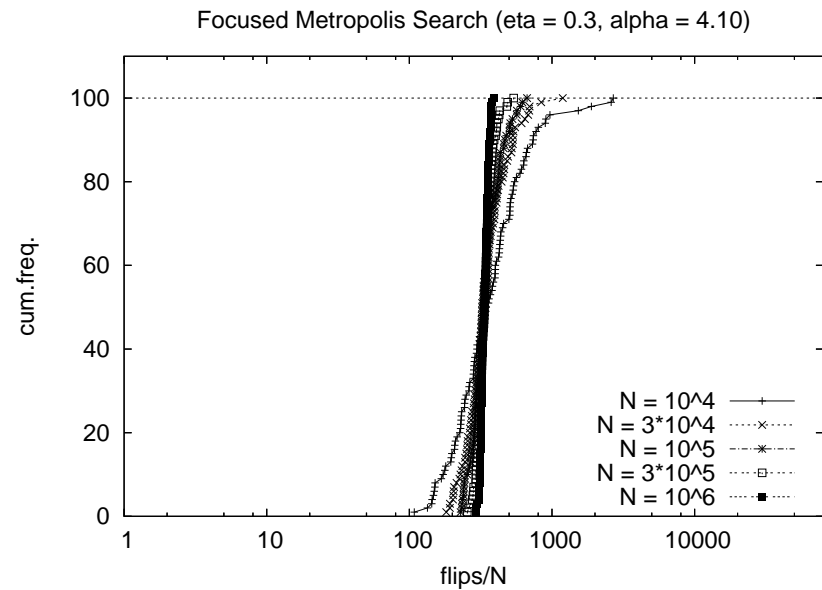
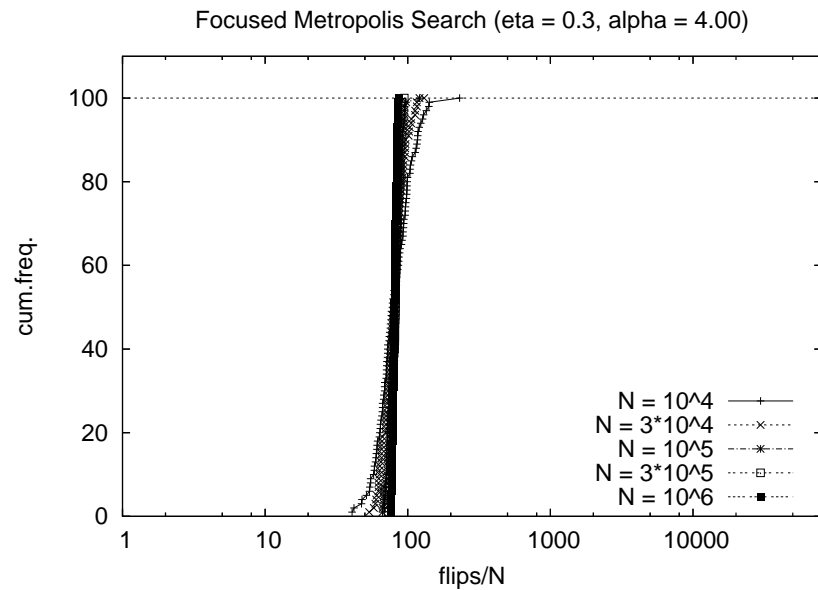
# FMS experiments (3-SAT)



Normalised solution times for FMS,  $\alpha = 3.8 \dots 4.3$ .  
Left: complete data; right: medians and quartiles.



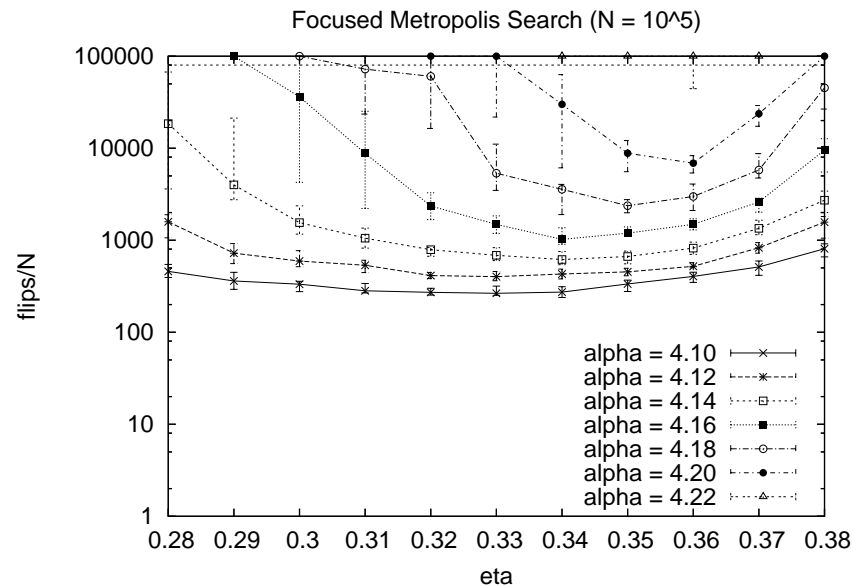
# FMS linear scaling



Cumulative solution time distributions for FMS with  $\eta = 0.3$ .



# FMS optimal acceptance ratio?

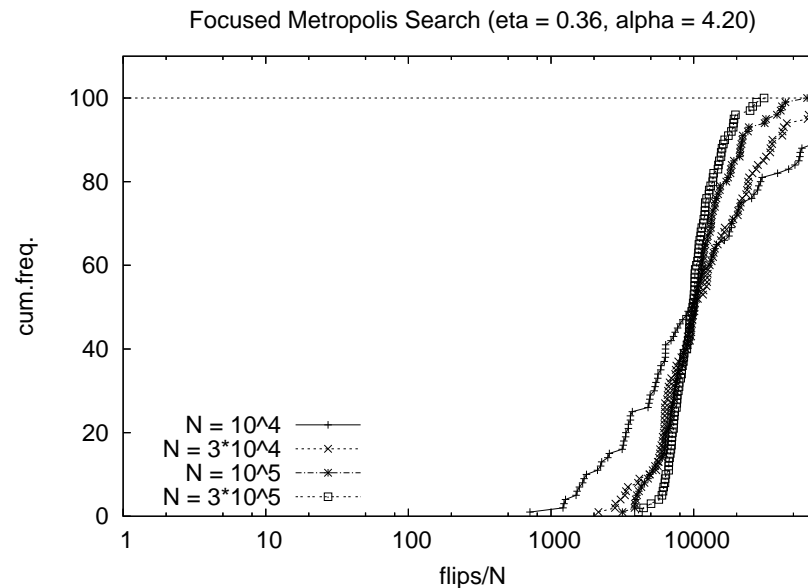


Normalised solution times for FMS with  $\eta = 0.28 \dots 0.38$ ,  
 $\alpha = 4.10 \dots 4.22$ .





# FMS optimal acceptance ratio cont'd



Cumulative solution time distributions for FMS with  $\eta = 0.36$ ,  
 $\alpha = 4.20$ .



# Analysis?

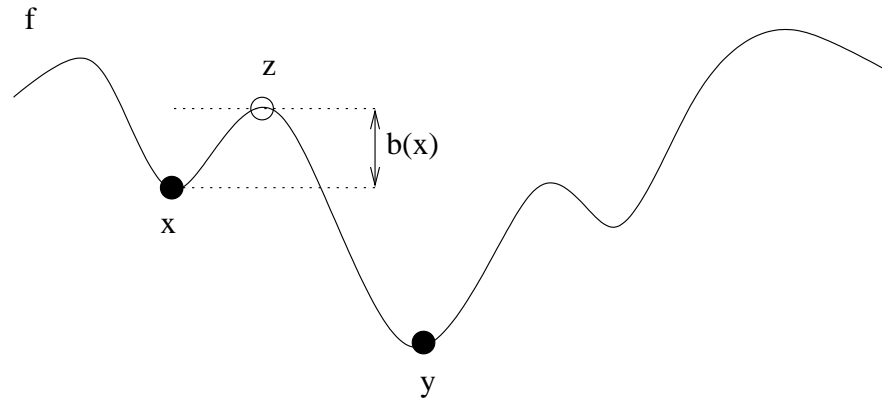
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- For FRRT: landscape structure?
- For FMS: contact processes?



# Combinatorial landscapes

Reidys & Stadler (SIAM Rev. 2002)

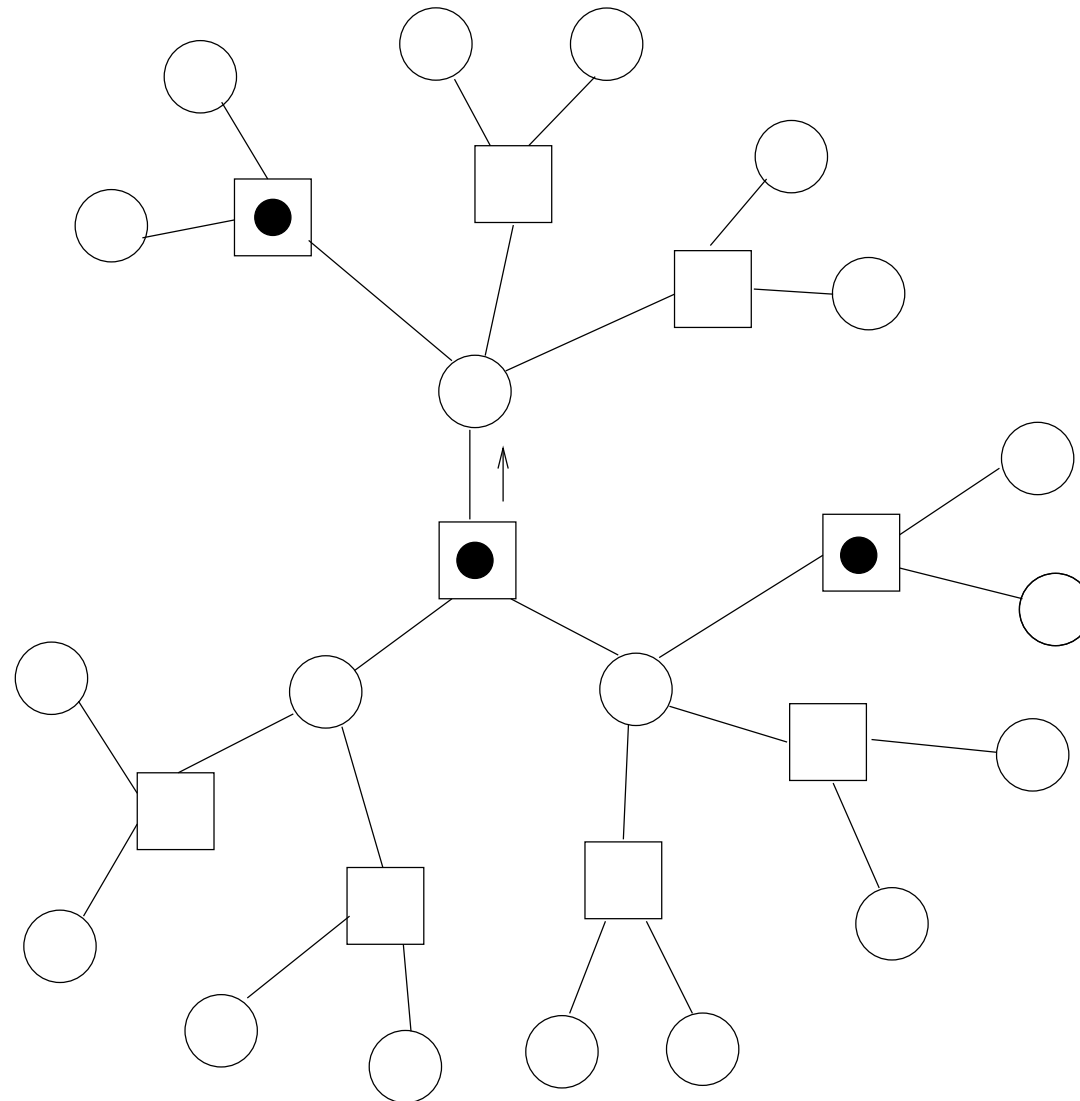


- *barrier height from x to y*:  $b(x, y) = \min\{\max\{f(z) - f(x), 0 \mid z \in p\} \mid p \text{ an } x\text{-}y \text{ path}\}$
- *barrier height of x*  $\notin \text{Opt}$ :  
 $b(x) = \min\{b(x, y) \mid f(y) < f(x)\}$
- *depth of a landscape*:  
 $D_f = \max\{b(x) \mid x \notin \text{Opt}\}$



# Focused search as a contact process

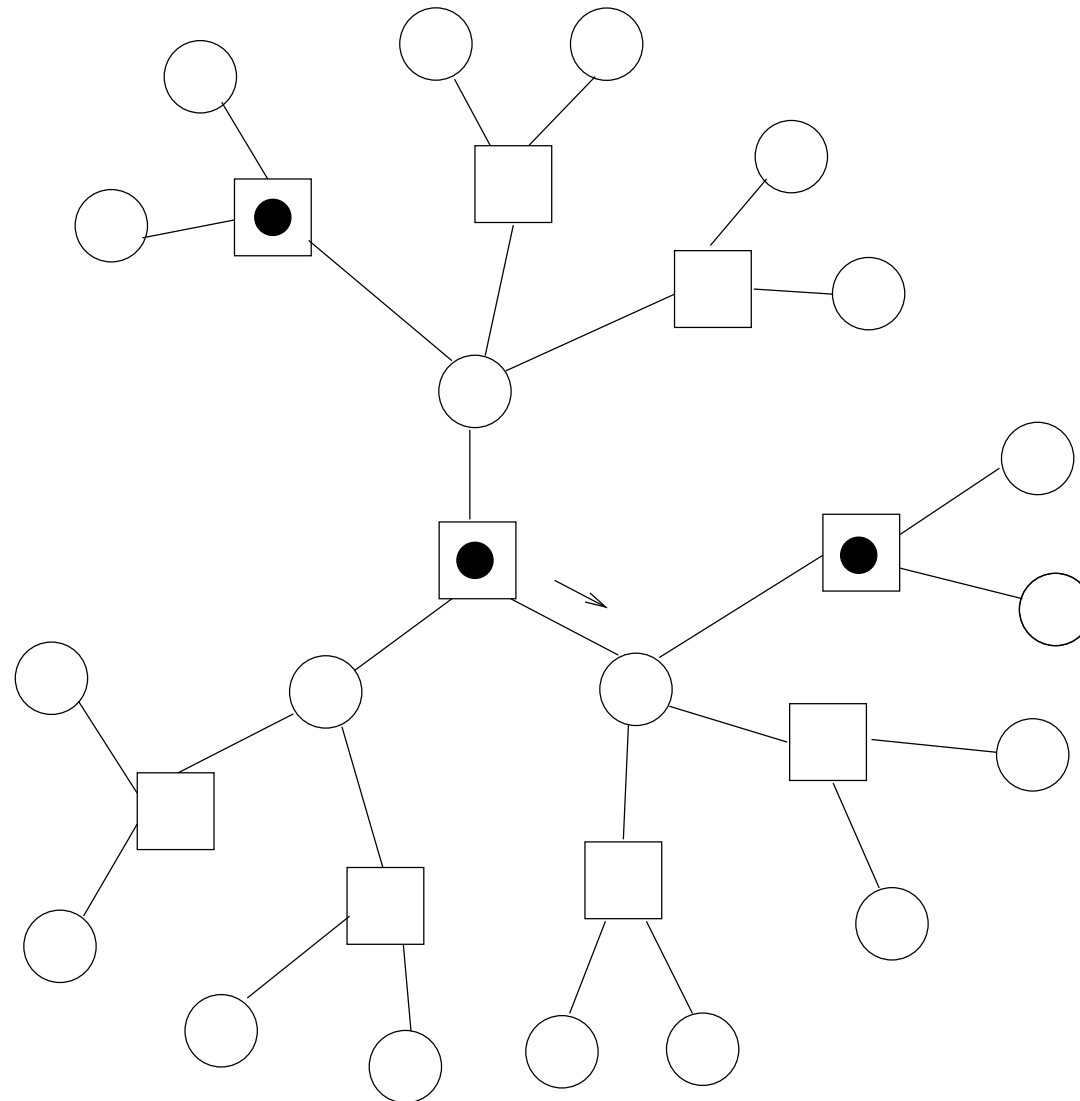
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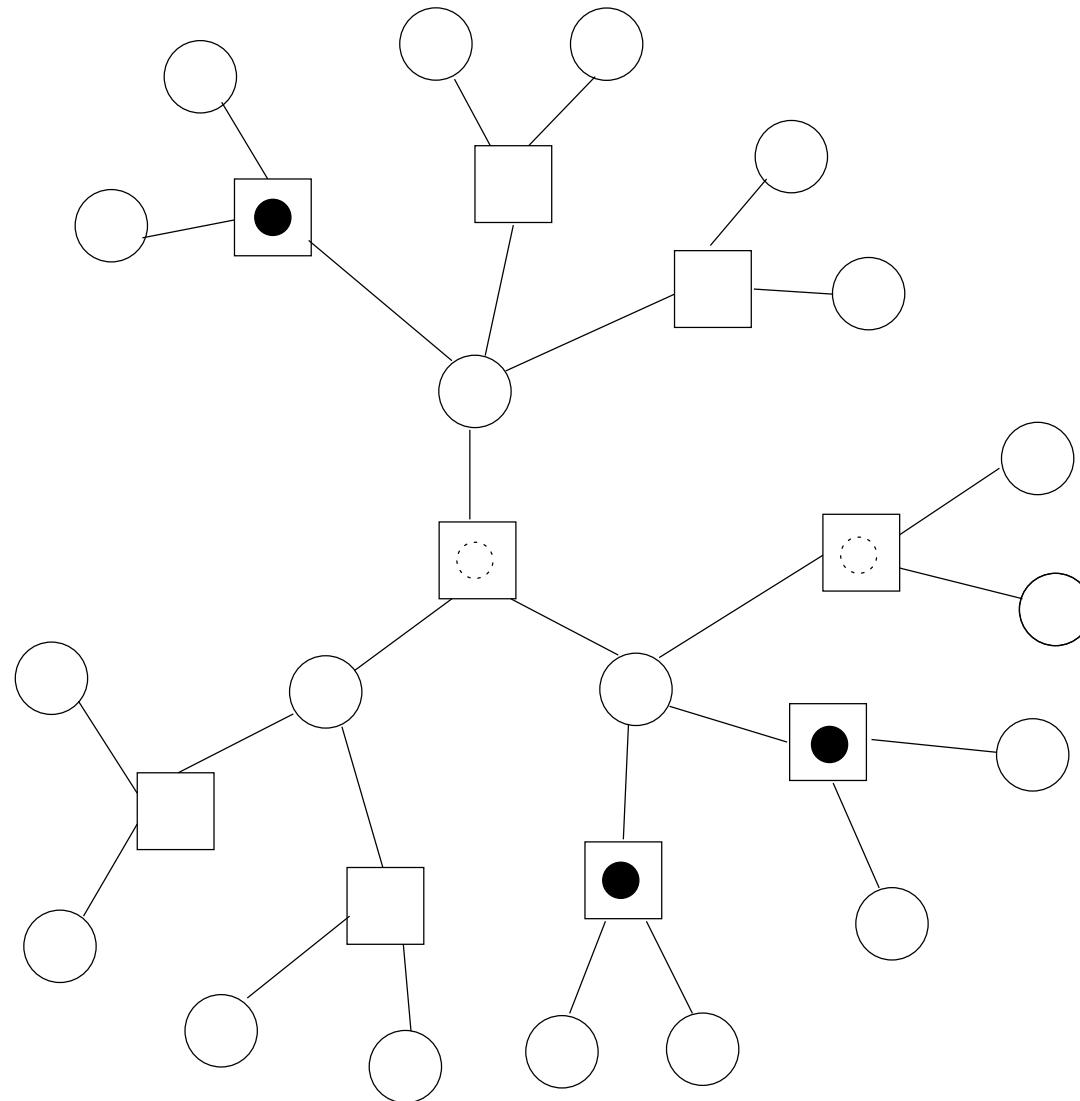
# Focused search as a contact process

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# Focused search as a contact process

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# References

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## References

- [1] S. Seitz, P. Orponen: An efficient local search method for random 3-satisfiability. *Proc. LICS'03 Workshop on Typical Case Complexity and Phase Transitions (Ottawa, Canada, June 2003)*. Elsevier Electronic Notes in Discrete Mathematics Vol. 16.
- [2] S. Seitz, M. Alava, P. Orponen: Threshold behaviour of WalkSAT and focused Metropolis search on random 3-satisfiability. *Proc. 8th Intl. Conf. on Theory and Applications of Satisfiability Testing (St. Andrews, Scotland, June 2005)*. Springer-Verlag, Berlin, to appear.
- [3] S. Seitz, M. Alava, P. Orponen: Focused local search algorithms for random 3-satisfiability. To appear.

