

## Combinatorial Models and Stochastic Algorithms

## Tutorial 9, April 1

## Problems

1. Compute the expected value of the clustering coefficient  $\mathcal{C}(G)$  for an ER random graph  $G \in \mathcal{G}(n, p)$ . Give also some estimates for the expected value of the characteristic path length  $\mathcal{L}(G)$ .
2. Compute the clustering coefficient  $\mathcal{C}(G)$ , characteristic path length  $\mathcal{L}(G)$ , and distribution of node degrees for a circulant graph  $C_{nk}$ . (It suffices to compute these quantities asymptotically for fixed  $k$  and large  $n$ .) What is the edge density  $p = e(C_{nk})/\binom{n}{2}$  for such a graph? What is the effect on  $\mathcal{L}(G)$  of a single randomly added shortcut edge?
3. Compute the clustering coefficient  $\mathcal{C}(G)$ , characteristic path length  $\mathcal{L}(G)$ , and distribution of node degrees for a “caveman graph” consisting of  $k$  “caves” of  $r$  nodes each. What is the edge density  $p$  for such a graph? (Recall that a “caveman graph” is a cyclic arrangement of  $k$  appropriately modified  $r$ -cliques. It suffices to compute these quantities asymptotically for (a) fixed  $r$  and large  $k$  and (b) the case  $r \sim k \sim \sqrt{n}$  for large  $n$ .)
4. Verify by direct calculation that in the simulated annealing algorithm the finite-temperature Gibbsian distributions  $\pi^{(T)}$  for  $T > 0$  do indeed converge pointwise to the desired limit distribution  $\pi^*$  as  $T \rightarrow 0$ .
5. Consider a simple state space graph with states  $S = \{\sigma_0, \sigma_1, \sigma_2, \sigma_3\}$  and neighbourhood structure  $N(\sigma_i) = \{\sigma_{i-1}, \sigma_{i+1}\}$ , where the indices are computed modulo 4. Write down explicitly the transition probability matrix of the simulated annealing algorithm at temperature  $t$  for this system, when the function to be minimised is given by  $H(\sigma_0) = 1$ ,  $H(\sigma_1) = 2$ ,  $H(\sigma_2) = 0$ ,  $H(\sigma_3) = 2$ . Given a cooling schedule where the temperature at step  $k$  of the algorithm is  $t_k > 0$ , what is the probability that the algorithm when initialised in the locally optimal state  $\sigma_0$  will stay there forever? Find a sequence  $t_k$  for which this probability is nonzero. What kind of cooling schedule would, according to Theorem 6.5 (p. 93 of the notes) guarantee asymptotic convergence to the globally optimal state  $\sigma_2$ ?
6. The NP-complete PARTITION problem is defined as follows: given a sequence of  $2n$  nonnegative integers  $x_1, \dots, x_{2n}$ , is there a subsequence of  $n$  numbers whose sum is exactly half the sum of the whole sequence? Formulate the task of finding approximate partitions of an integer sequence as a minimisation problem, and present a simulated annealing approach to solving it. What kinds of cooling schedules would Theorem 6.5 suggest for your algorithm in the case of input sequences consisting of  $2n$  numbers from the interval  $[0, N]$ ? (You might consider also actually implementing your algorithm and experimenting with some more realistic cooling schedules.)