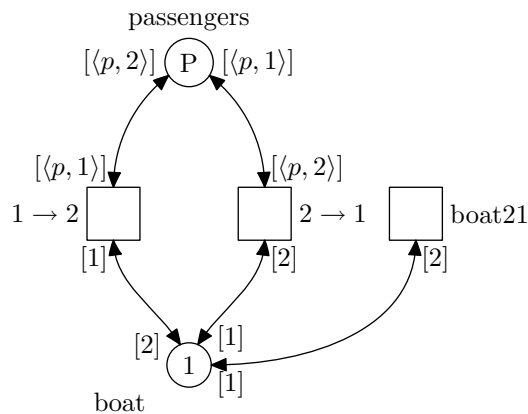


1. For each LTL formula shown below, give a finite automaton where the formula holds in the initial state and an automaton where the formula does not hold. State the which propositions hold in which states.

- (a) $\Box(p \rightarrow \Diamond q)$
- (b) $(p \mathbf{U} q) \vee (\Box \neg q)$
- (c) $\Box \Diamond p$
- (d) $\Diamond p \rightarrow (\neg p \mathbf{U} q)$

2. Consider the following algebraic net modelling the ferryman's problem. Present LTL formulae for the following properties. Do the formulae hold in the model? If not, give a counterexample.



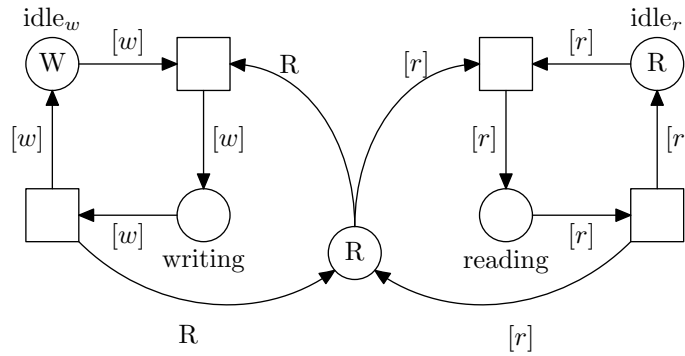
- (a) All passengers reach the opposite shore
- (b) The wolf eats the goat
- (c) It is not possible that any of the passengers will be eaten

3. Solve the following problem with modulo 22 arithmetics:

Traverse all numbers $0 \leq i \leq 21$ by doing 22 additions in modulo 22 arithmetics. The numbers in the additions are restricted in such a way that only some numbers are allowed, and even then only a restricted number of times. The numbers allowed are: 18 three times, 17 five times, 14 four times, 8 three times, 6 three times, and 5 four times. Start the additions from 0.

- (a) Design an algebraic net to solve the problem
- (b) Write a Maria description of the net
- (c) Use a reject statement to find a solution
- (d) From a computational point of view, was this a sensible approach to solving the problem?

4. A demonstration whose parts b) and c) are beyond the scope of the course. Below is presented a net modelling a reader/writer problem. Let $W = [1]$ ja $R = [1, 2]$.



- (a) Perform reachability analysis for the net
- (b) Construct a Büchi automaton corresponding to the reachability graph with atomic propositions are $p = \text{“reader 1 is reading”}$ and $q = \text{“reader 2 is reading”}$
- (c) Model check the formula $\diamond(p \wedge q)$.