

Homework problems:

1. Let $\Sigma = \{a, b\}$. Give some examples of strings from each of the following languages (at least three strings per language):
 - (a) $\{a^n b^{2n} \mid n \geq 0\}$;
 - (b) $\{w \in \Sigma^* \mid w \text{ contains exactly two occurrences of the substrings } ba \text{ and/or } ab\}$;
 - (c) $\{w \in \Sigma^* \mid \exists u, v \in \Sigma^* \text{ s.e. } w = uvv\}$;
 - (d) $\{w \in \Sigma^* \mid w = w^R\}$
2. The *reversal* of a string $w \in \Sigma^*$, denoted w^R , is defined inductively by the rules:
 - (i) $\varepsilon^R = \varepsilon$;
 - (ii) if $w = ua$, where $u \in \Sigma^*$ and $a \in \Sigma$, then $w^R = aw^R$.

It was proved in class (cf. also Lewis & Papadimitriou, p. 43) that for any strings $u, v \in \Sigma^*$ it is the case that $(uv)^R = v^R u^R$. Prove in a similar manner, by induction based on the above definition of reversal, the following facts:

- (a) $(w^R)^R = w$;
 - (b) $(w^k)^R = (w^R)^k$, for any $k \geq 0$.
3. Prove that if A and B are countably infinite sets, then the union $A \cup B$ and the intersection $A \cap B$ are both countable.

Demonstration problems:

4. Show that any alphabet Σ with at least two symbols is comparable to the binary alphabet $\Gamma = \{0, 1\}$, in the sense that strings over Σ can be easily encoded into strings over Γ and vice versa. How much can the length of a string change in your encoding? (I.e., if the length of a string $w \in \Sigma^*$ is $|w| = n$ symbols, what is the length of the corresponding string $w' \in \Gamma^*$?) Could you design a similar encoding if the target alphabet consisted of only *one* symbol, e.g. $\Gamma = \{1\}$?
5. Prove that the Cartesian product $\mathbb{N} \times \mathbb{N}$ is countably infinite. (*Hint:* Think of the pairs $(m, n) \in \mathbb{N} \times \mathbb{N}$ as embedded in the Euclidean (x, y) plane \mathbb{R}^2 . Enumerate the pairs by diagonals parallel to the line $y = -x$.) Conclude from this result that also the set \mathbb{Q} of rational numbers is countably infinite.
6. Let S be an arbitrary nonempty set.
 - (a) Give some injective (i.e. one-to-one) function $f : S \rightarrow \mathcal{P}(S)$.
 - (b) Prove that there cannot exist an injective function $g : \mathcal{P}(S) \rightarrow S$. (*Hint:* Assume that such a function g existed. Consider the set $R = \{s \in S \mid s \notin g^{-1}(s)\}$, and denote $r = g(R)$. Is it then the case that $r \in R$?)

Observe, as a consequence of item (b), that the power set $\mathcal{P}(S)$ of any countably infinite set S is uncountable.