

Homework problems:

1. Design a right-linear grammar that generates the language

$$\{w \in \{a, b\}^* \mid w \text{ does not contain substring } bba\}$$

(Cf. Demonstration Problem 3/1b.)

2. (a) Show that the following context-free grammar is ambiguous:

$$\begin{aligned} S &\rightarrow ASb \mid A \mid b \\ A &\rightarrow aA \mid a \end{aligned}$$

- (b) Design an unambiguous grammar for the same language. Describe the language informally.

3. Construct a context-free grammar for the language:

$$\{a^i b^j a^k \mid 0 \leq i \leq j \text{ or } i = k\}$$

Is your grammar ambiguous?

Demonstration problems:

4. Prove that the class of context-free languages is closed under unions, concatenations, and the Kleene star operation, i.e. if the languages $L_1, L_2 \subseteq \Sigma^*$ are context-free, then so are the languages $L_1 \cup L_2$, $L_1 L_2$ and L_1^* .
5. (a) Prove that the following context-free grammar is ambiguous:

$$\begin{aligned} S &\rightarrow \mathbf{if } b \mathbf{ then } S \\ S &\rightarrow \mathbf{if } b \mathbf{ then } S \mathbf{ else } S \\ S &\rightarrow s. \end{aligned}$$

- (b) Design an unambiguous grammar that is equivalent to the grammar in item (a), i.e. that generates the same language. (*Hint*: Introduce new nonterminals B and U that generate, respectively, only “balanced” and “unbalanced” **if-then-else**-sequences.)
6. Design a recursive-descent (top-down) parser for the grammar from Problem 6/6.