Helsinki University of Technology, Laboratory for Theoretical Computer Science TJ T-79.144 Logic in Computer Science: Foundations Examination, October 4, 2004

Please note the following: your answers will be graded only if you have passed all the three home assignments before the exam!

Assignment 1 Answer and justify exactly (at most half a page per item).

- (a) True or false: it is possible to define the other propositional connectives  $(\neg, \land, \lor, \leftrightarrow)$  using the connectives  $\rightarrow$  and  $\vee$  (exclusive or).
- (b) True or false: if Σ<sub>1</sub> and Σ<sub>2</sub> are sets of sentences such that Σ<sub>1</sub> ⊆ Σ<sub>2</sub> and φ is a sentence such that Σ<sub>1</sub> ⊨ φ, then also Σ<sub>2</sub> ⊨ φ.
- (c) True or false: a conjunctive normal form  $\phi$  of a sentence in predicate logic is logically equivalent to the form  $\phi'$  obtained from  $\phi$  by Skolemization.
- (d) True or false: the satisfiability problem SAT of propositional logic is **NP**-complete.

**Assignment 2** Examine if the given claim holds using semantic tableaux. If not, justify by giving a valuation/structure (a counter example).

- (a)  $\models \neg (A \land \neg B) \land (\neg C \to A) \to (A \land B) \lor (\neg A \land C)$
- (b)  $\{\exists x \exists y P(x,y), \forall x \forall y (P(x,y) \rightarrow Q(x,y))\} \models \exists x Q(x,x)$
- (c)  $\{\forall x \neg (A(x) \leftrightarrow B(x)), \forall y A(y) \lor \forall y \neg A(y)\} \models \forall z B(z) \lor \forall z \neg B(z)$

Tableau proofs must contain all intermediary steps !!!

## Assignment 3

(a) Derive a clausal form for the sentence

$$\neg (\neg \exists y E(y) \to \forall y (\exists x E(x) \to E(y))).$$

Try to make it as simple as possible.

(b) Consider the following program P:

z=0; v=x; while(!(z==y)) {z=z+1; v=v-1}

Use weakest preconditions and a suitable invariant to establish

 $\models_p [true] P [v == x - y].$ 

Assignment 4 Let us represent natural numbers 0, 1, 2, ... with ground terms 0, s(0), s(s(0)), ... built of a constant symbol 0 and a function symbol s which is interpreted as the function s(x) = x + 1 for natural numbers x.

- (a) Let the predicates J2(x), J3(x) and J6(x) mean that a natural number x is divisible by two, three and six, respectively. Use predicate logic to define these predicates such that the definition of the predicate J6 is based on the definitions of the predicates J2 and J3.
- (b) Use resolution to show that if a natural number x is divisible by two and three, then the natural number x + 6 is divisible by six.

The name of the course, the course code, the date, your name, your student id, and your signature must appear on every sheet of your answers.

Feedback: http://www.tcs.hut.fi/Studies/T-79.144/feedback.html