## Helsinki University of Technology Laboratory for Theoretical Computer Science Harri Haanpää (puh. 5243)

## T-79.1001 Introduction to Theoretical Computer Science T (4 cr) Exam 20 August 2007 noon to 3 p.m.

Write on every answer sheet:

- Name, degree programme, student number
- The text: "T-79.1001 Introduction to Theoretical Computer Science T, 20 August 2007"
- The total number of answer sheets submitted for grading
  - 1. Finite state automata and regular expressions.
    - a) Design a deterministic finite state automaton that recognizes the language  $L = \{w \in \{a,b\}^* \mid w \text{ has an odd number of both } as \text{ and } bs\}$  5 *p*.
    - b) Design a regular expression that describes the language  $L = \{w \in \{a, b\}^* \mid w \text{ has an odd number of } as \text{ or } bs \text{ (or both)}\}$ 5 *p*.
    - c) Find the minimal deterministic finite state automaton that accepts the language  $L = \{w \in \{a, b\}^* \mid w \text{ has an odd number of either } as \text{ or } bs (but not both)\}$  5 *p*.
  - 2. a) Design a context-free grammar for the language

$$L = \{ucvcw \mid u, v, w \in \{0, 1\}^*, v = u^R \text{ or } v = w^R \text{ (or both)} \}.$$

(Notation  $x^R$  denotes the reverse of string *x*, i.e. string *x* written backwards.) 5 *p*.

- b) Show that the grammar you gave in part (a) is ambiguous.
- c) Prove (precisely!) that the language in part (a) is not regular. (*Hint:* Consider e.g. strings of the form a<sup>n</sup>ca<sup>n</sup>cb<sup>n</sup>.) 5 p.
- Design a single-tape Turing machine that decides whether the input is of the form wcw, where w ∈ {a,b}\*. Present the computation of your machine with inputs abab and abcab.
- 4. a) Closure properties of language classes. Show that if the language *L* is regular, then the language  $\overline{L}$  (the complement of *L*) is also regular. 5*p*.
  - b) Show that if the languages  $L_1$  and  $L_2$  are context-free, then the language  $L_1 \cup L_2$  is also context-free. 5*p*.
  - c) Show that if the languages  $L_1$  and  $L_2$  are recursively enumerable, then the language  $L = L_1 \cap L_2$  is also recursively enumerable. 5*p*.

Total 60p.

5 p.